



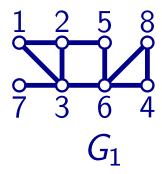
# Simultaneous Drawing of Planar Graphs with Right-Angle Crossings and Few Bends

Alexander Wolff
Chair of Computer Science I
Universität Würzburg

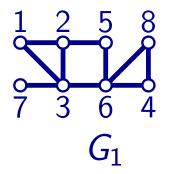
Joint work with Michael A. Bekos · Thomas C. van Dijk · *Philipp Kindermann* 

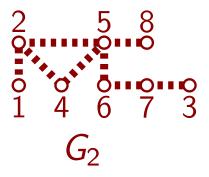
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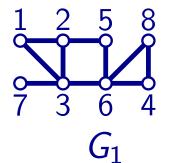


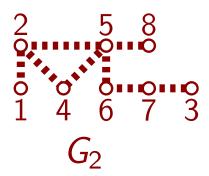
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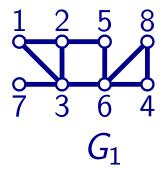


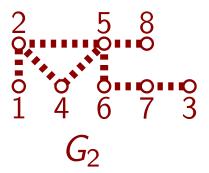


Embed both graphs in a planar way

edges of one graph may intersect edges of the other graph

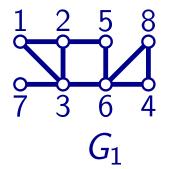
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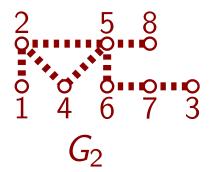




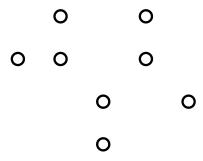
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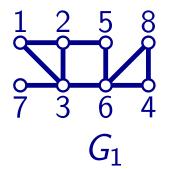


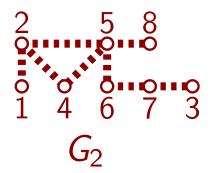


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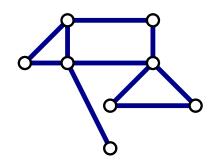


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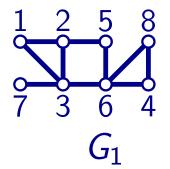


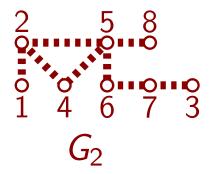


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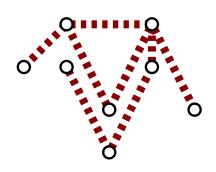


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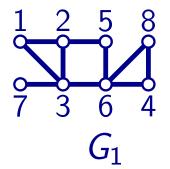


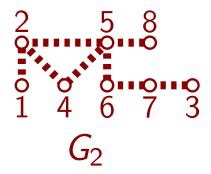


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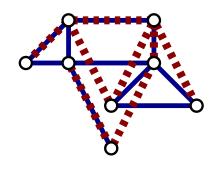


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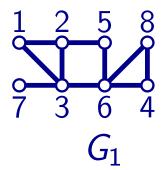


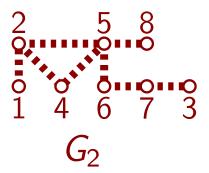


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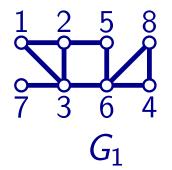
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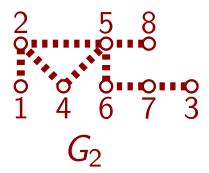




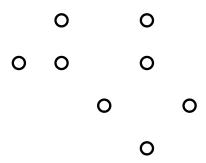
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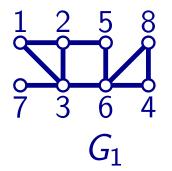


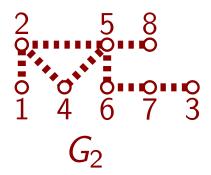


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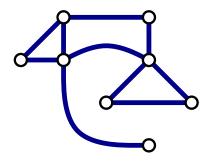


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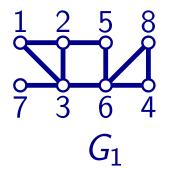


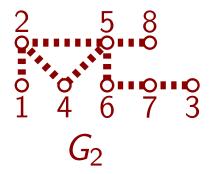


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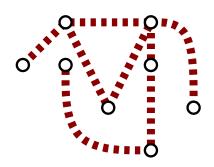


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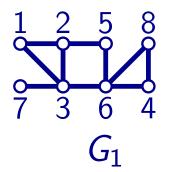


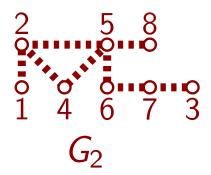


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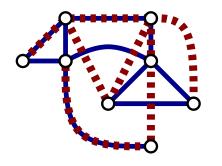


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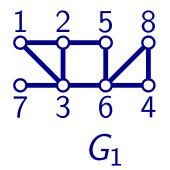


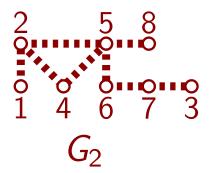


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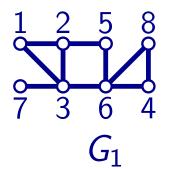
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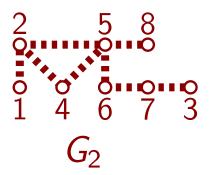




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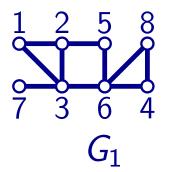
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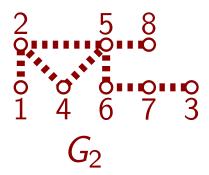




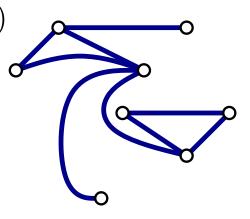
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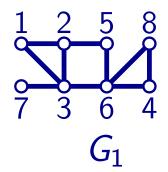


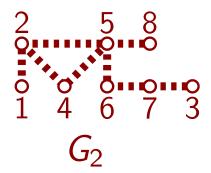


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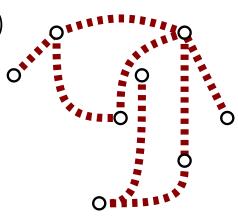


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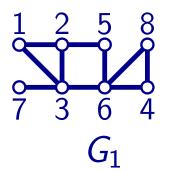


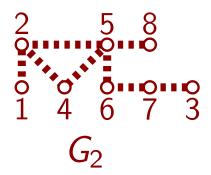


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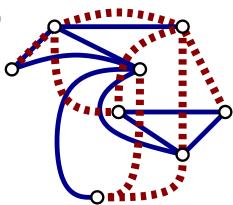


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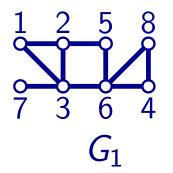


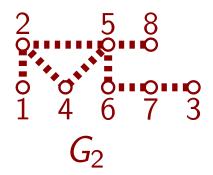


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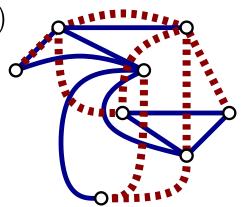




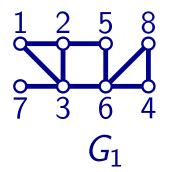
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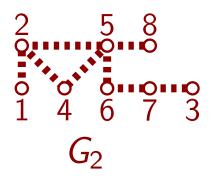
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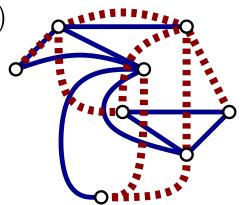




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RACSIM: Simultaneous Embedding with Right-Angle Crossings

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Geometric RACSIM: RACSIM with straight-line edges

#### [Angelini et al. JGAA'12]

There exist a tree and a path that don't admit an  ${\rm SGE}$  drawing.

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 $\circ$  A tree and a planar graph admit a (6, 8)-SEFE.

#### A Few Known Results

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Two trees admit a (1, 4)-SEFE
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New!

- $\circ$  A tree and a planar graph admit a (6, 8)-SEFE.
- Two planar graphs admit a (6, 16)-SEFE.

## Our Results for SE

#### Graph classes

#### Number of bends

Cycle	X	Cycle	$1 \times 1$
Caterpillar	X	Cycle	1  imes 1
Four Matchings			1  imes 1  imes 1  imes 1
Tree	×	Matching	$1 \times 0$
Wheel	X	Matching	$2 \times 0$
Outerpath	×	Matching	$2 \times 1$
Outerplanar	X	Outerplanar	$3 \times 3$
2-page book emb.	X	2-page book emb.	$4 \times 4$
Planar	×	Planar	$6 \times 6$

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Graph classes		Number of bends	
Cycle	×	Cycle	$1 \times 1$
Caterpillar	×	Cycle	1  imes 1
Four Matchings			1  imes 1  imes 1  imes 1
Tree	×	Matching	1  imes 0
Wheel	×	Matching	$2 \times 0$
Outerpath	×	Matching	$2 \times 1$
Outerplanar	×	Outerplanar	$3 \times 3$
2-page book emb.	×	2-page book emb.	$4 \times 4$
Planar	×	Planar	$6 \times 6$

Bend complexity can be seen as a measure that shows how difficult it is to simultaneously embed two graphs.

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#### Graph classes

#### Number of bends

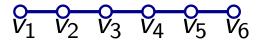
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Four Matchings			1  imes 1  imes 1  imes 1
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Outerpath	×	Matching	$2 \times 1$
Outerplanar	×	Outerplanar	$3 \times 3$
2-page book emb.	×	2-page book emb.	$4 \times 4$
Planar	X	Planar	$6 \times 6$

Bend complexity can be seen as a measure that shows how difficult it is to simultaneously embed two graphs.



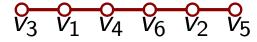


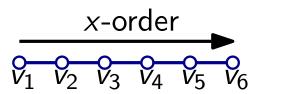


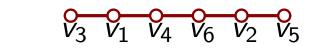


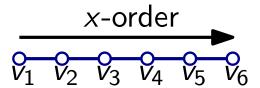


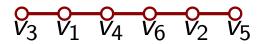


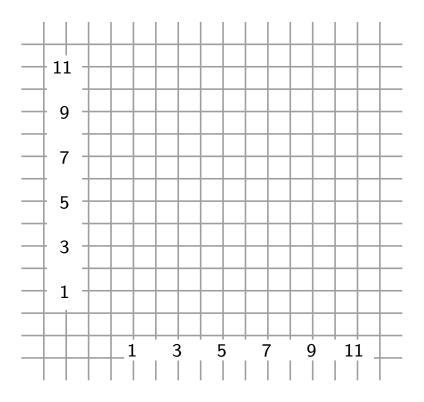


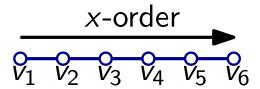


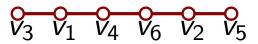


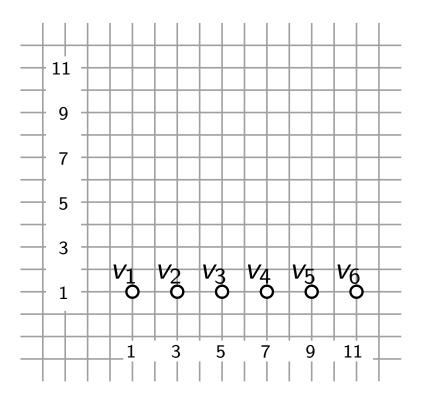


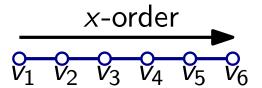




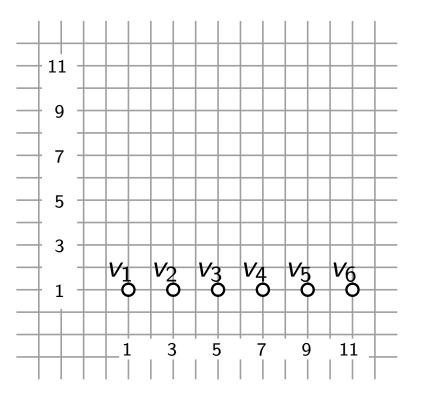


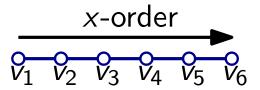


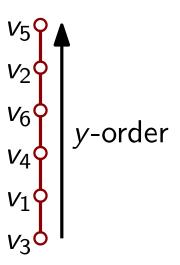


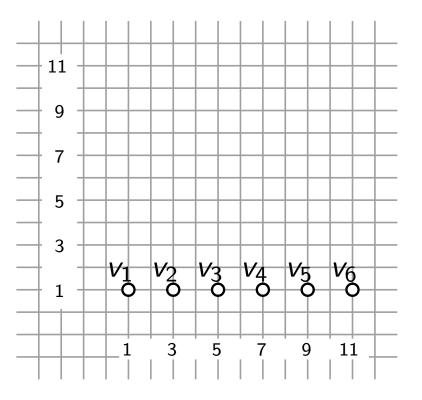


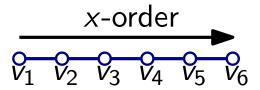


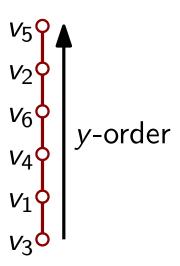


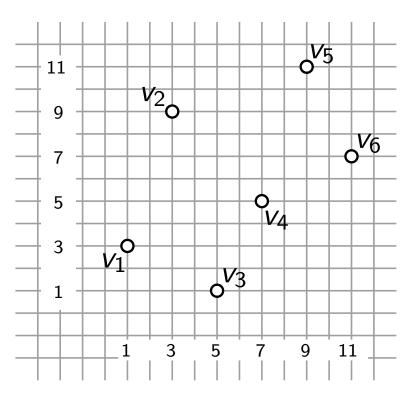


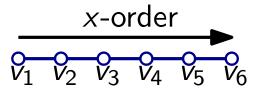


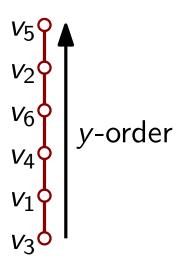




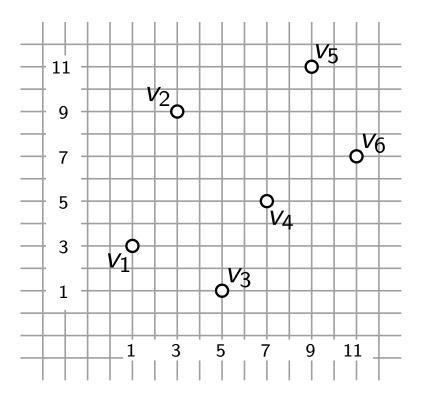


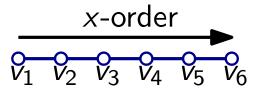


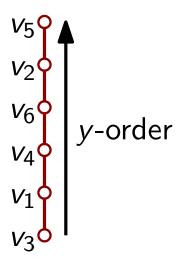


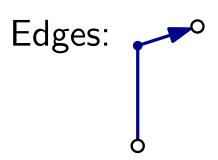


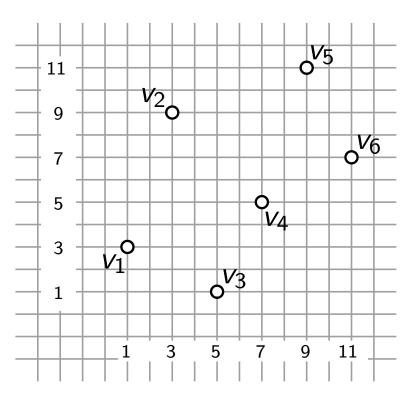
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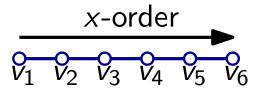


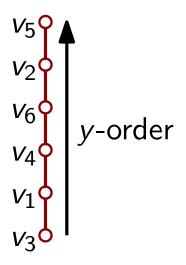


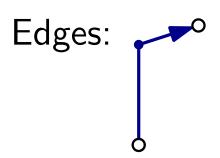


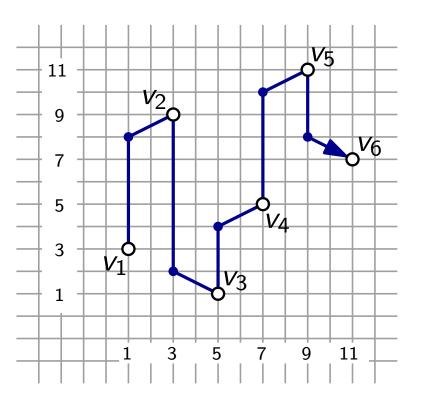


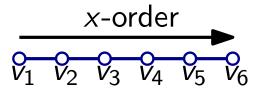


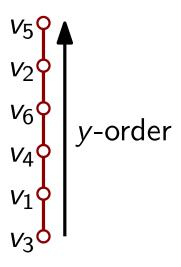


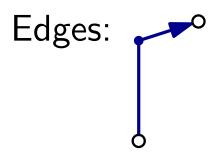




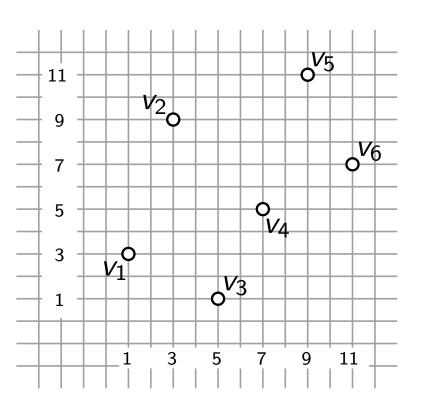


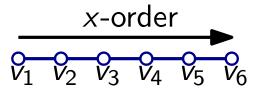


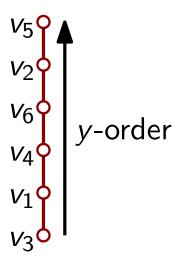


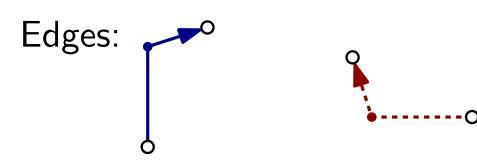


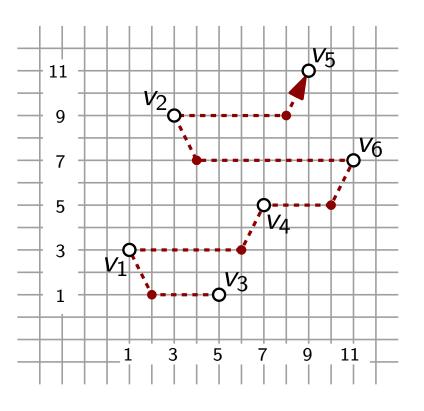


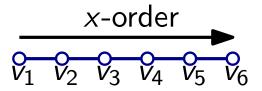


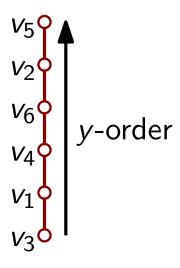


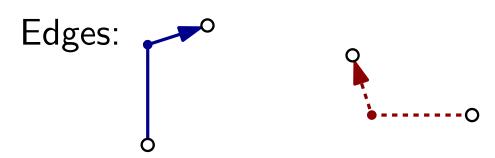


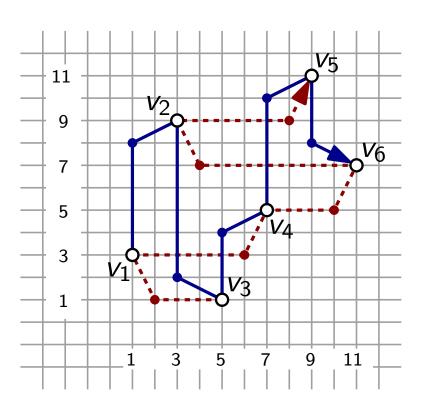


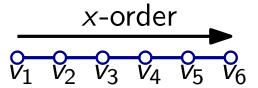


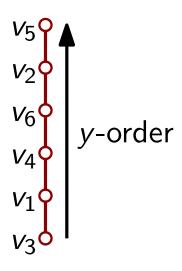


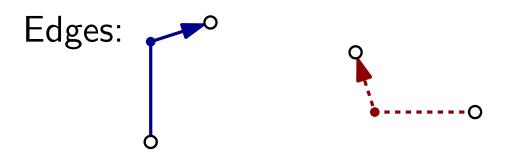






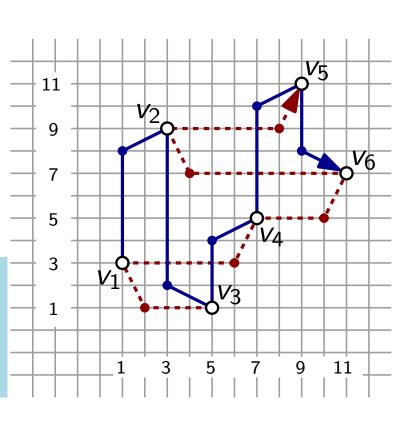


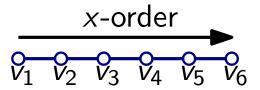


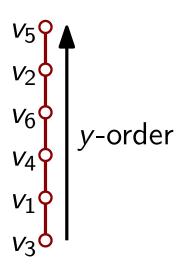


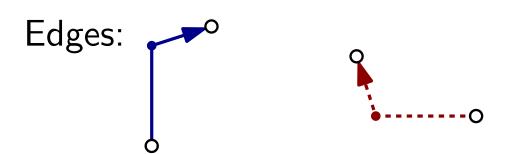
#### Main ideas:

Combine x-order and y-order.

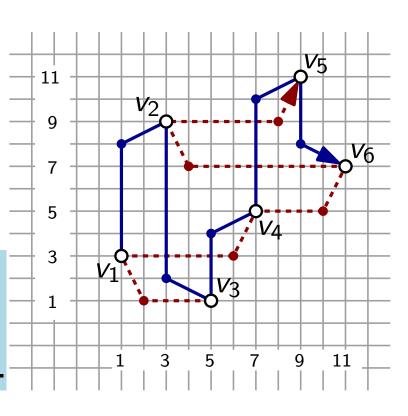


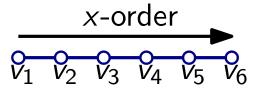


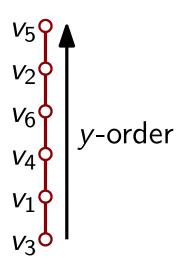


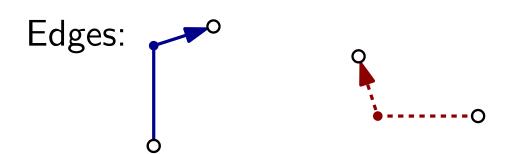


- Combine x-order and y-order.
- Keep slanted segm. short in 1 dim.

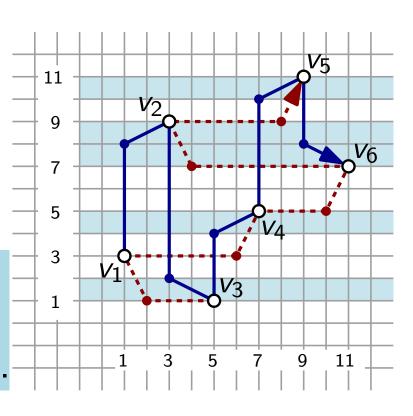


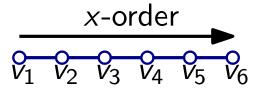


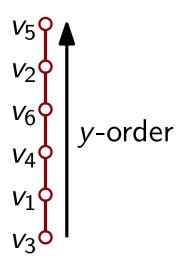


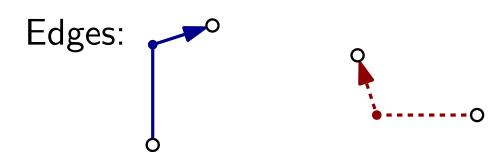


- Combine x-order and y-order.
- Keep slanted segm. short in 1 dim.

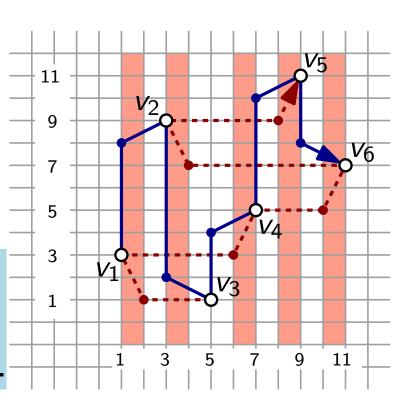


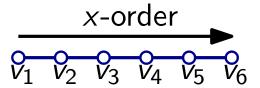


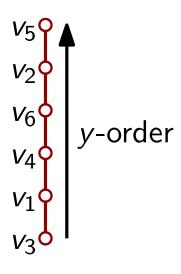


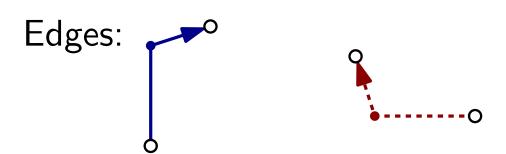


- Combine x-order and y-order.
- Keep slanted segm. short in 1 dim.

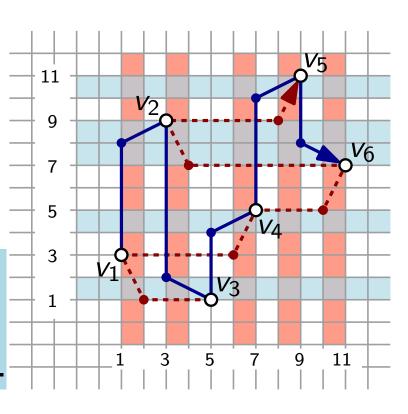




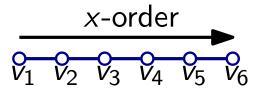


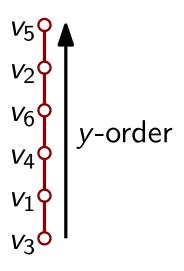


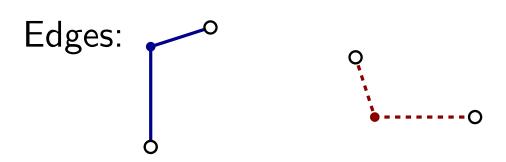
- Combine x-order and y-order.
- Keep slanted segm. short in 1 dim.

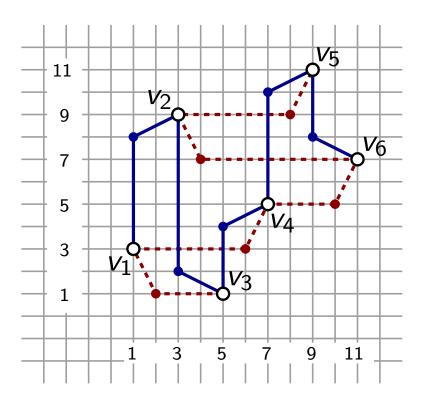


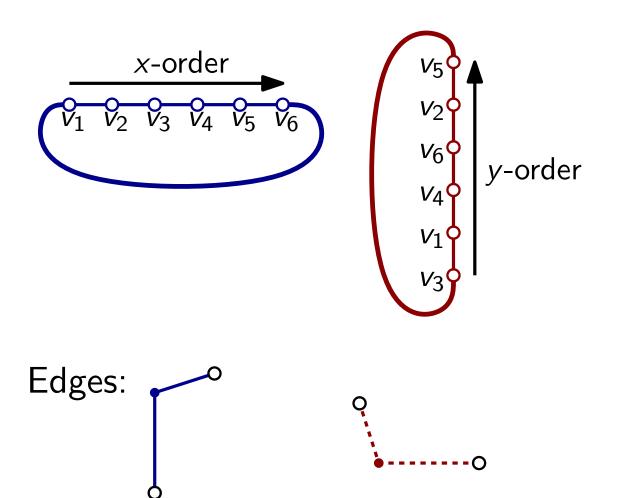
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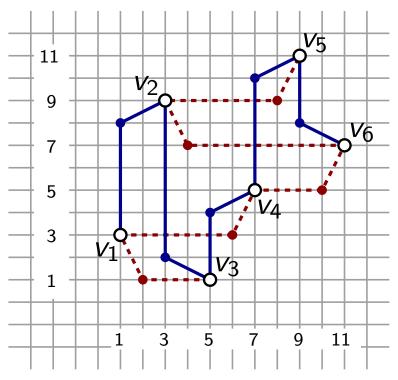


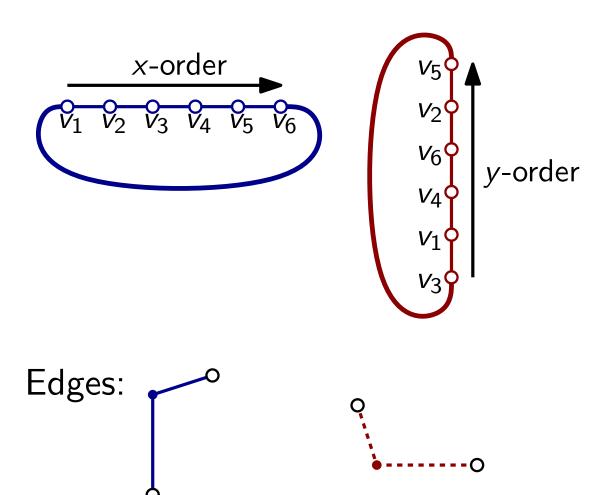


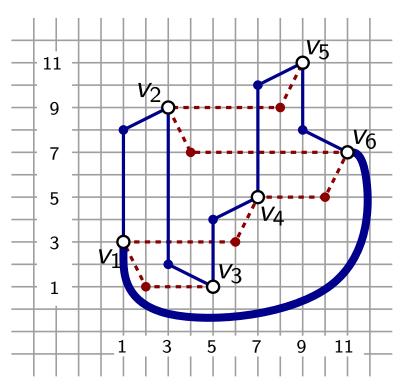


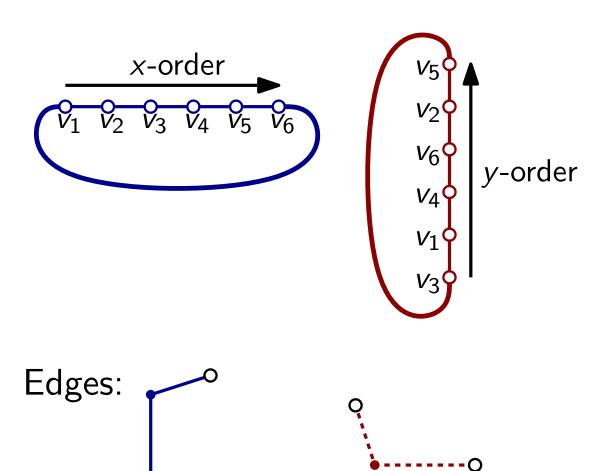


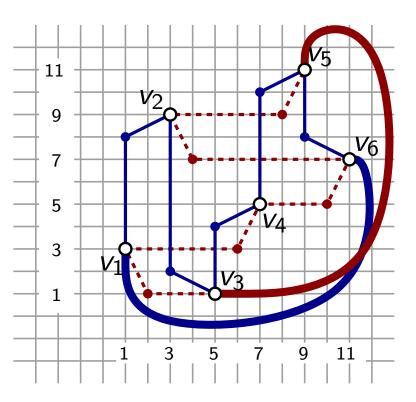




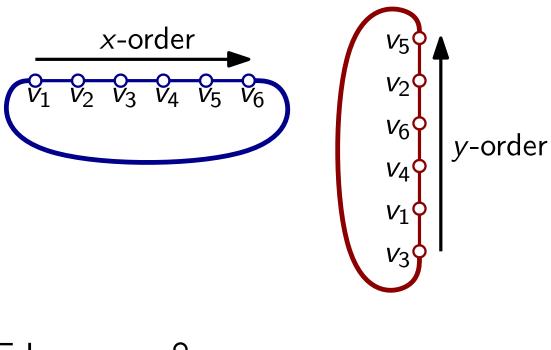


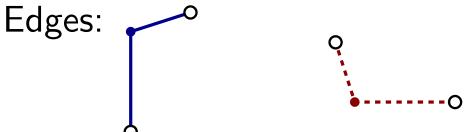


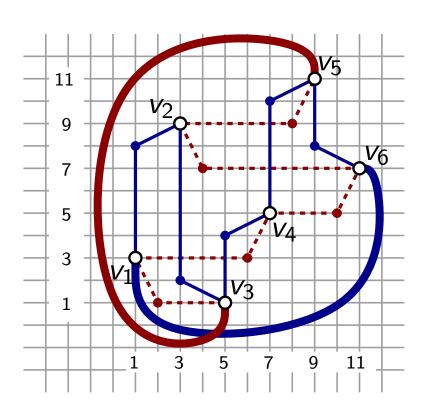


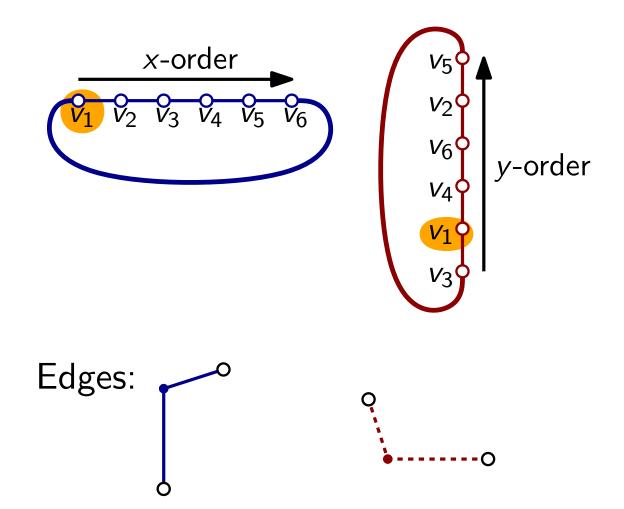


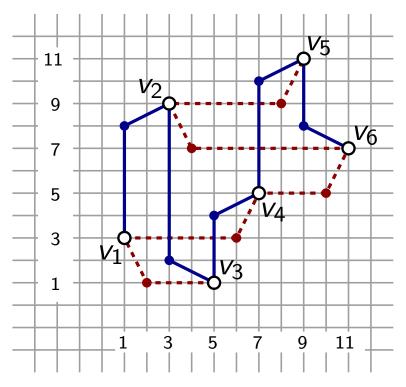
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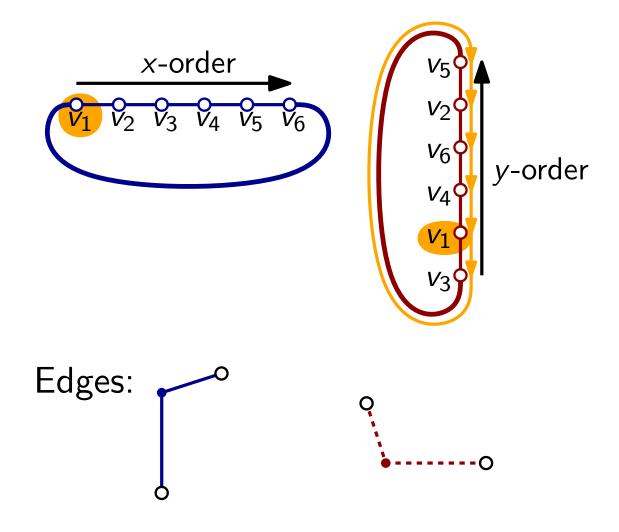


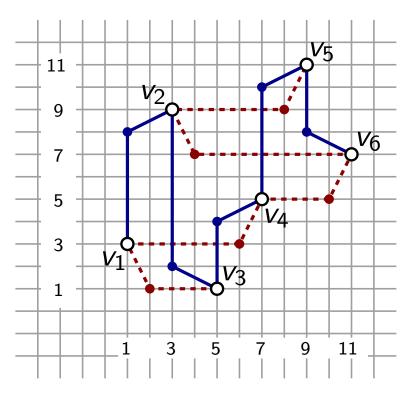


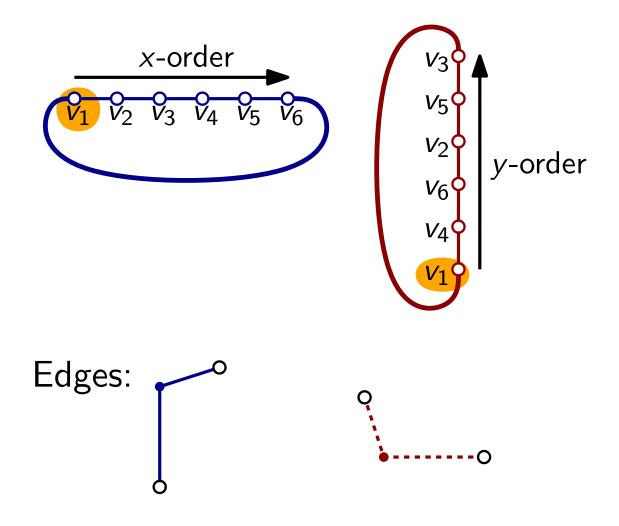


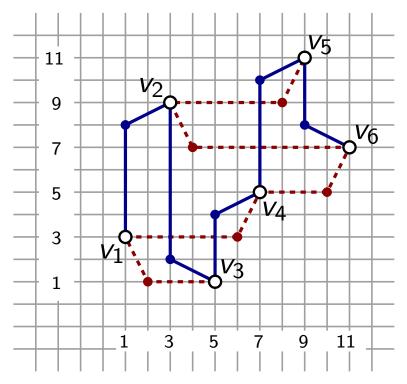


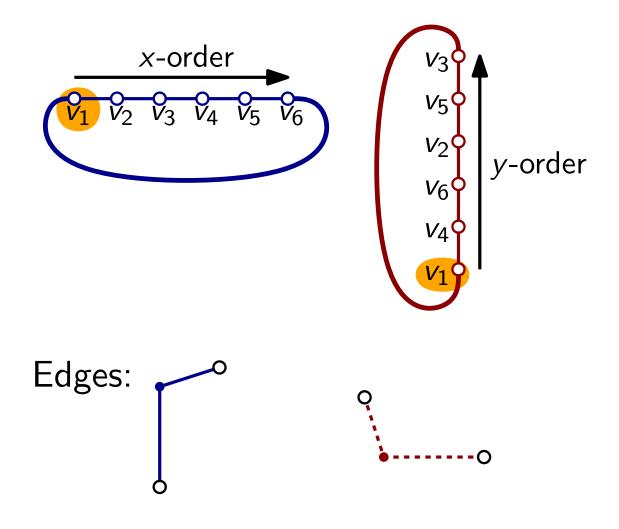
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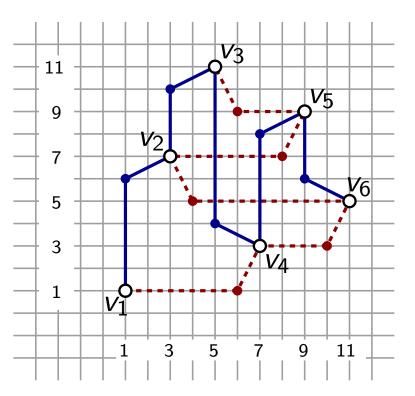


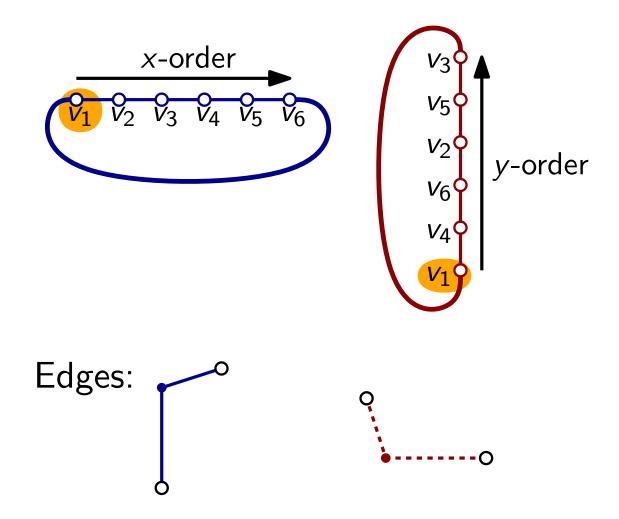


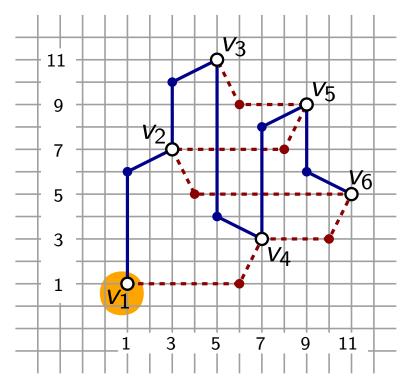


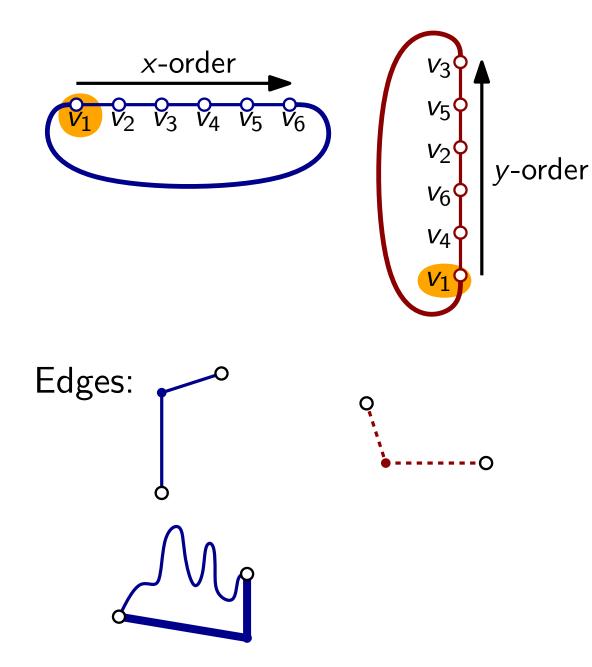


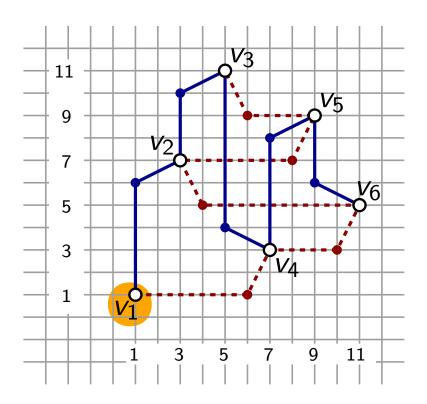




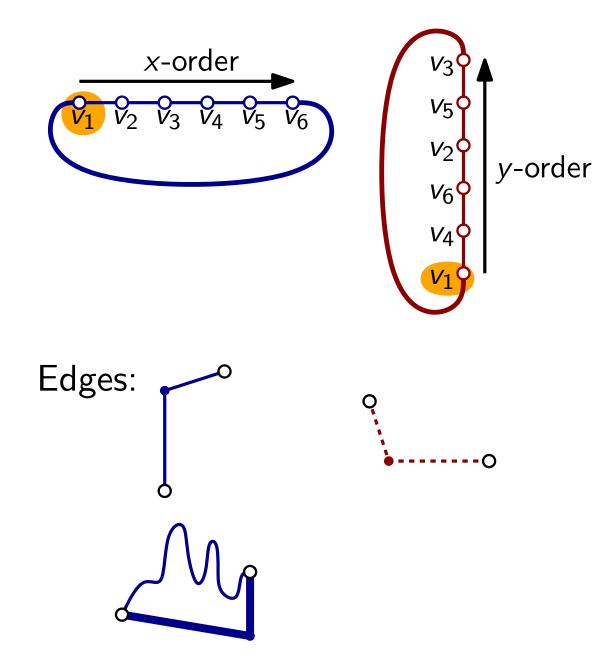


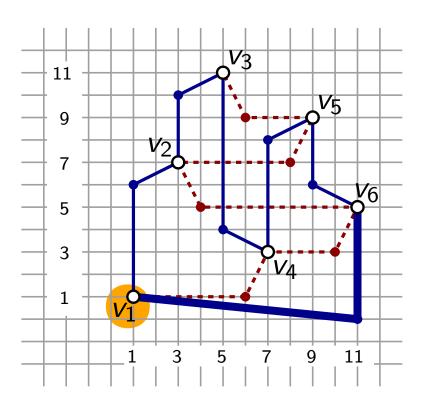




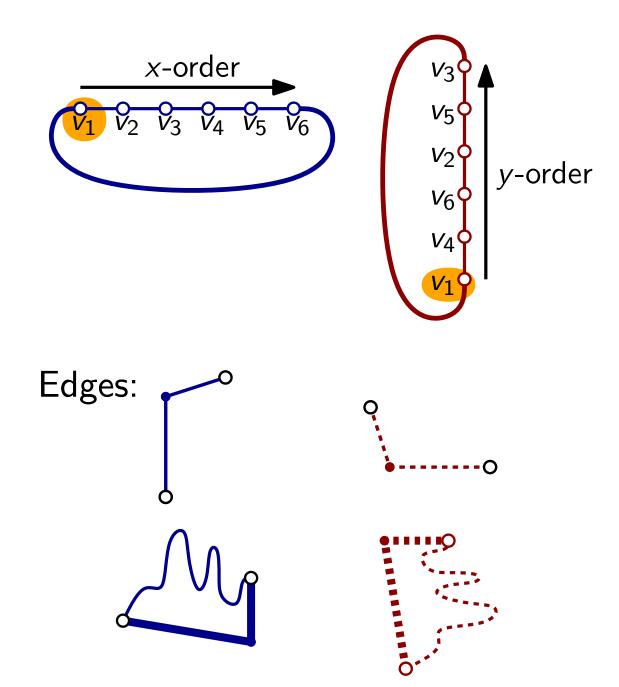


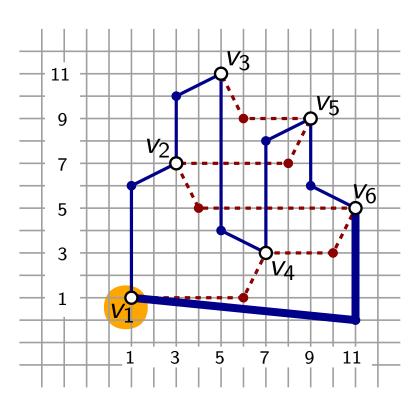
## Cycle × Cycle



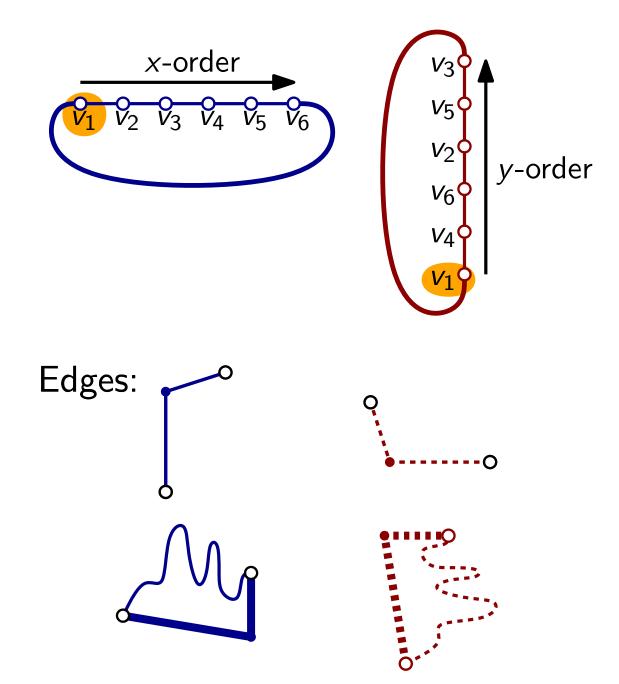


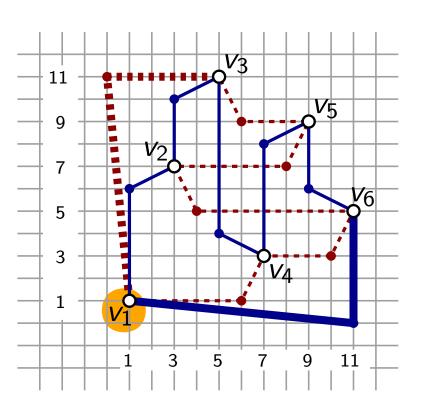
# Cycle $\times$ Cycle



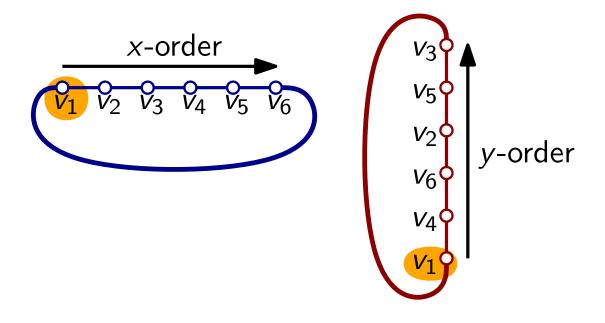


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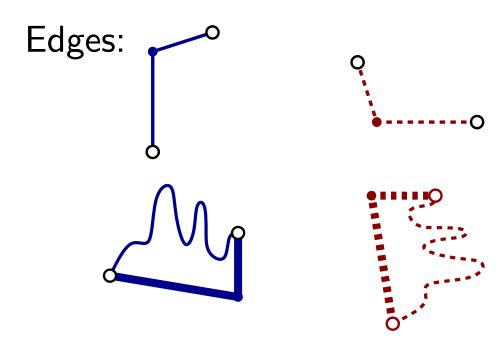


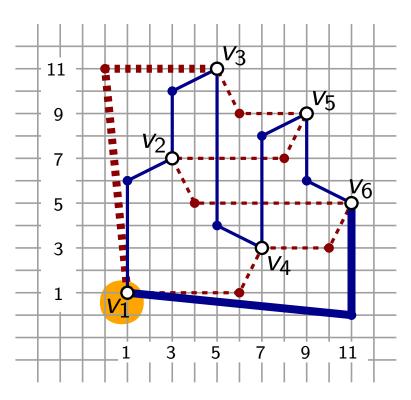


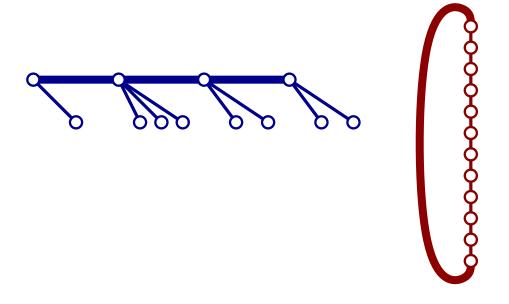
## Cycle × Cycle



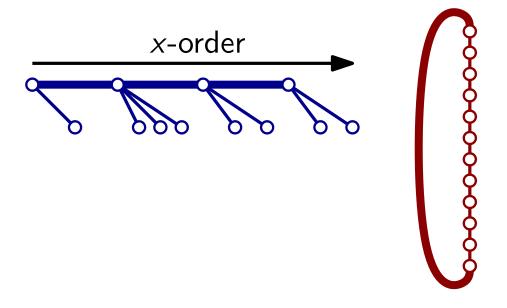
Bends:  $1 \times 1$ Grid size:  $(2n-1)^2$ 

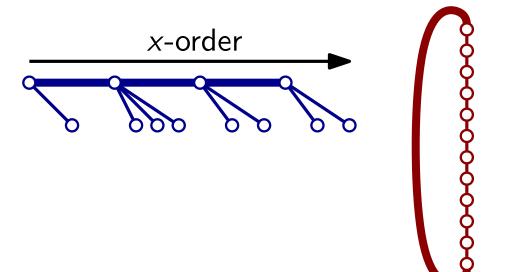


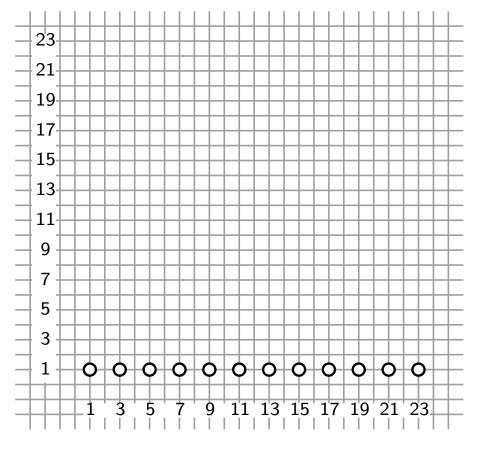


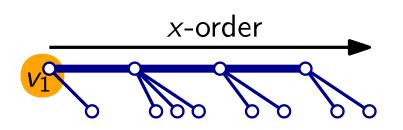


# $\mathsf{Caterpillar} \times \mathsf{Cycle}$

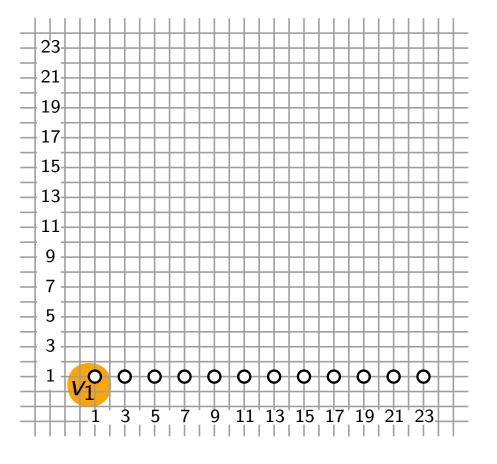




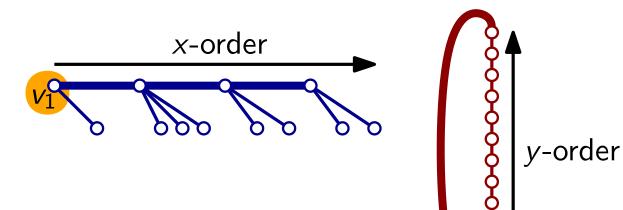


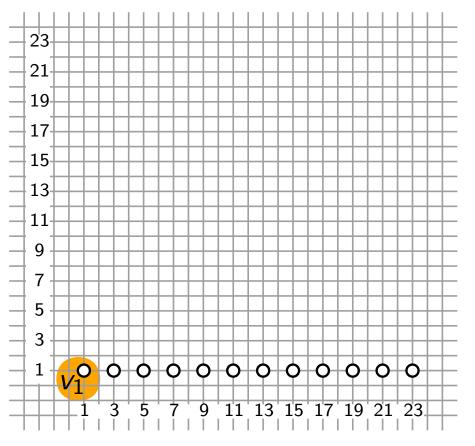




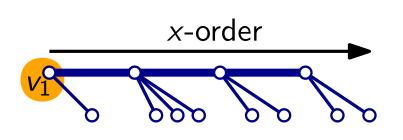


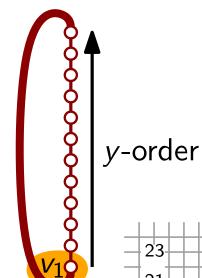
# Caterpillar × Cycle

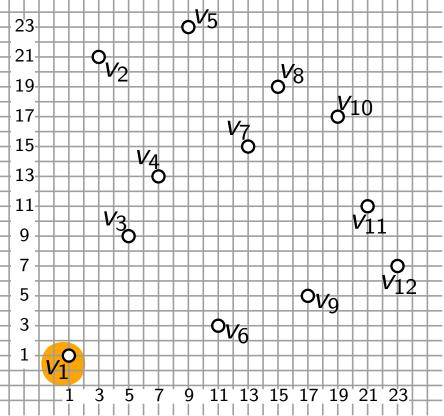




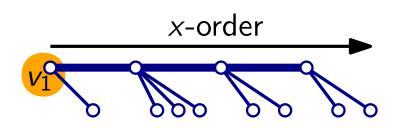
# Caterpillar × Cycle



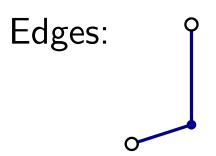


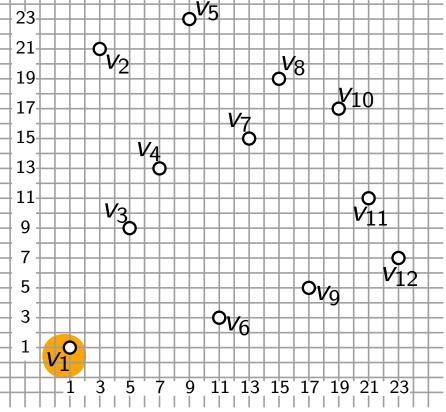


# ${\sf Caterpillar} \times {\sf Cycle}$

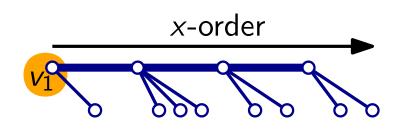


y-order

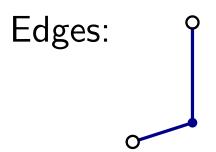


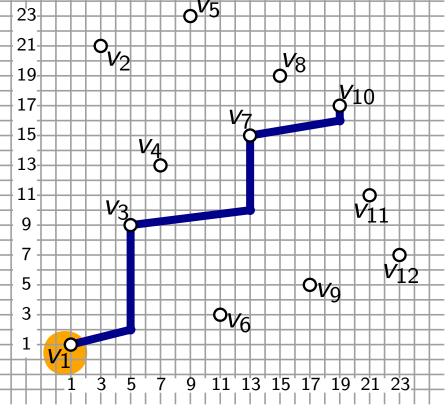


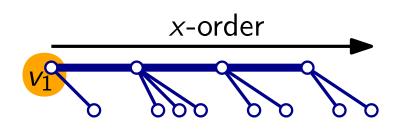
# ${\sf Caterpillar} \times {\sf Cycle}$



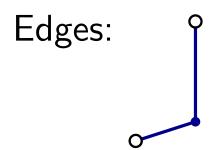
y-order

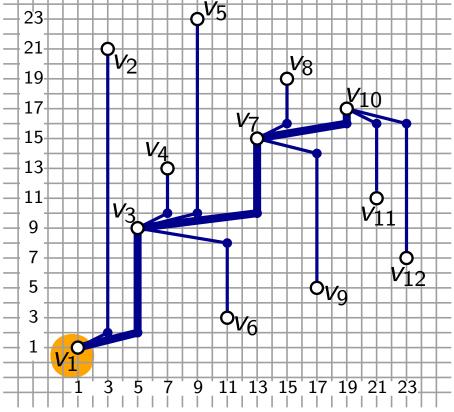


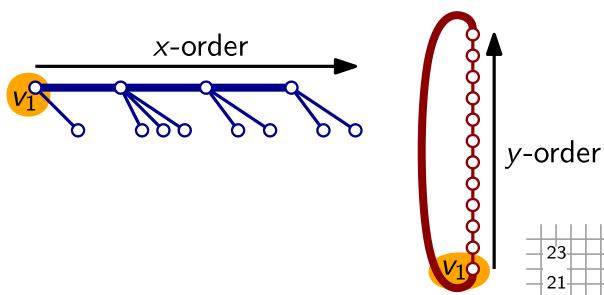


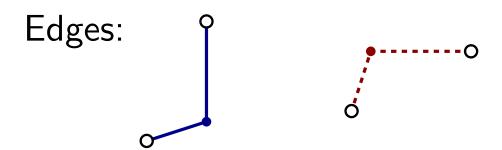


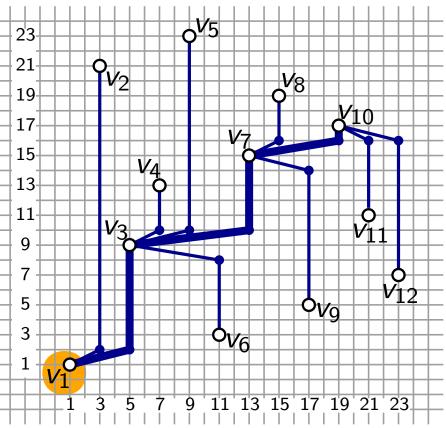
y-order

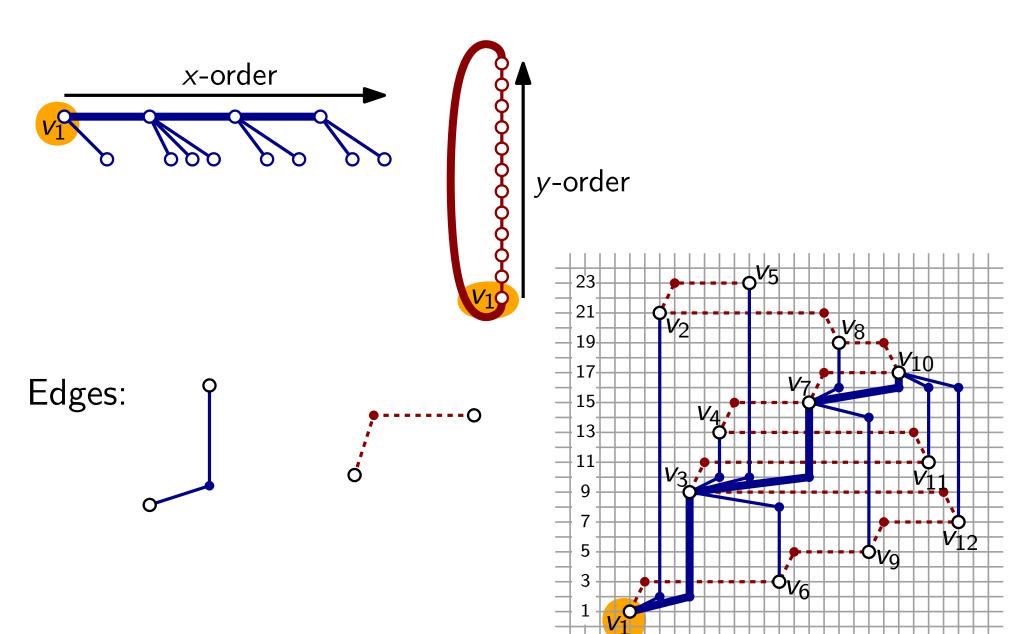


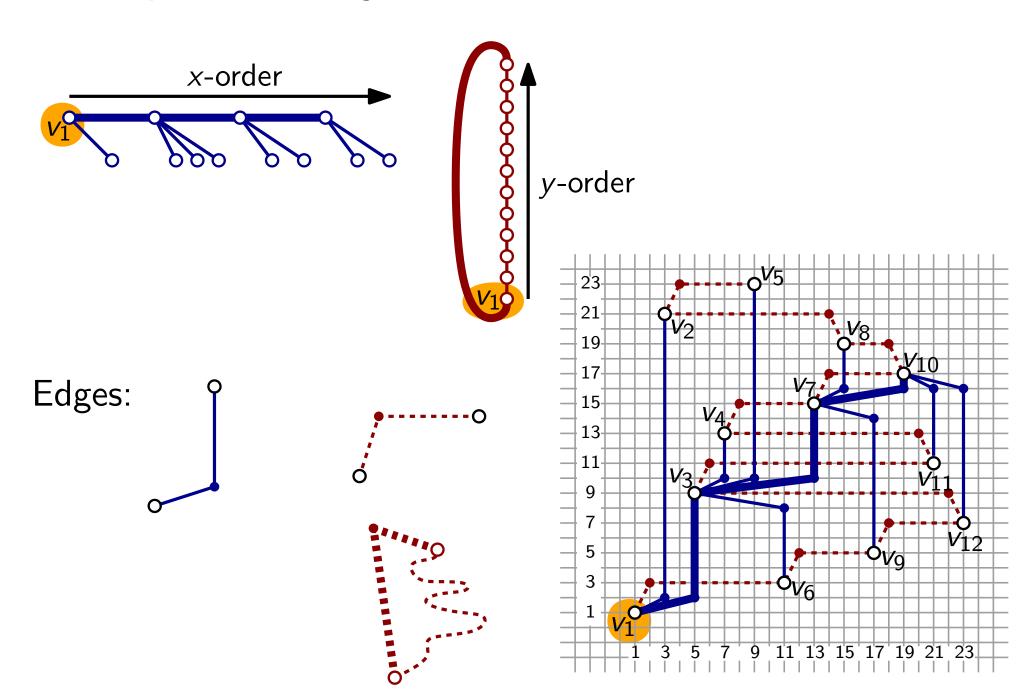


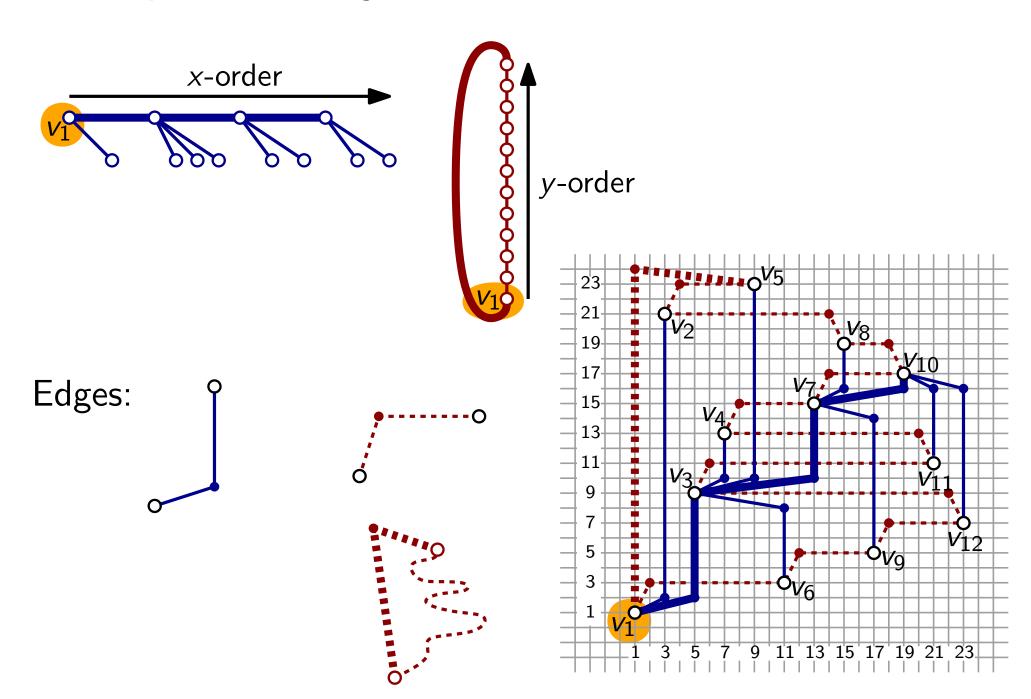


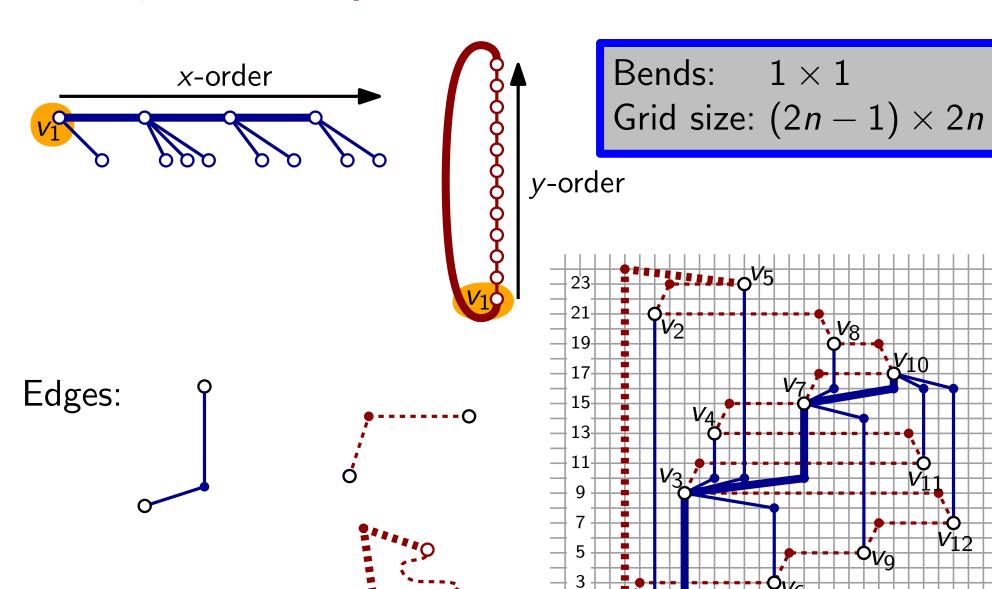










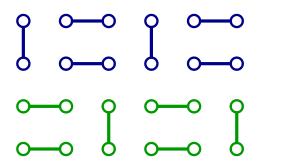


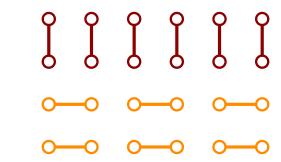
## Overview

### Graph classes

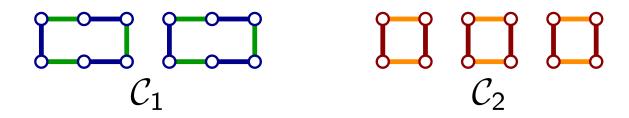
#### Number of bends

Cycle	×	Cycle	$1 \times 1$
Caterpillar	×	Cycle	1  imes 1
Four Matchings			1  imes 1  imes 1  imes 1
Tree	×	Matching	$1 \times 0$
Wheel	×	Matching	$2 \times 0$
Outerpath	×	Matching	$2 \times 1$
Outerplanar	×	Outerplanar	$3 \times 3$
2-page book emb.	X	2-page book emb.	$4 \times 4$
Planar	X	Planar	$6 \times 6$



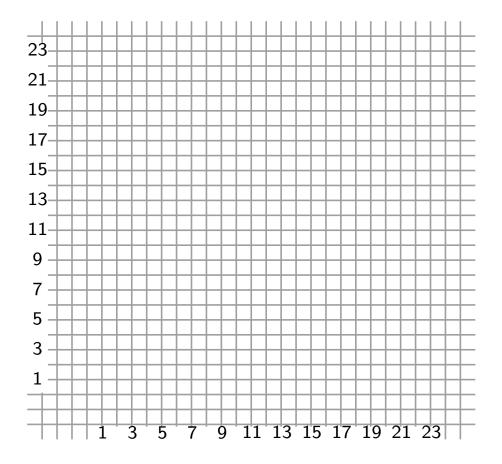


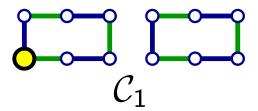
Combine  $\Rightarrow$  two sets of cycles  $C_1$ ,  $C_2$ 

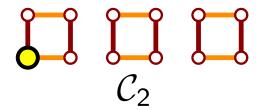


Combine  $\Rightarrow$  two sets of cycles  $C_1$ ,  $C_2$ 



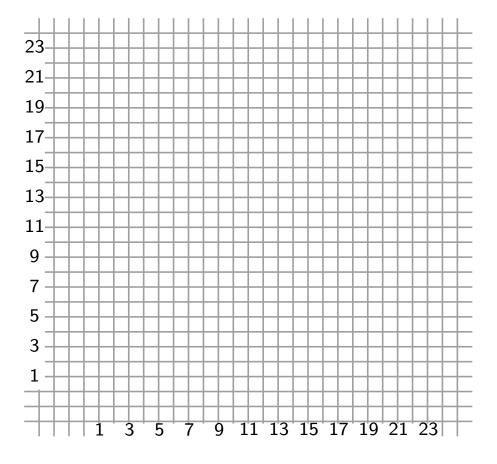


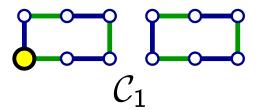


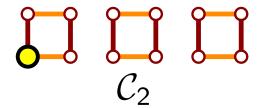


### Placement algorithm:

 $\circ$  Pick  $v_1$ 

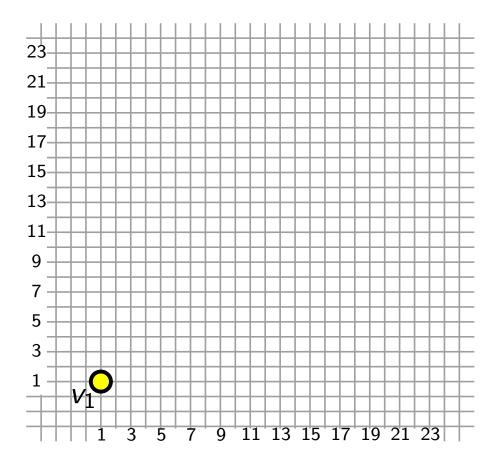


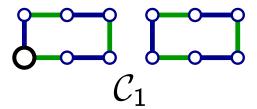


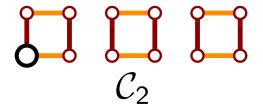


### Placement algorithm:

 $\bigcirc$  Pick  $v_1$ , place it

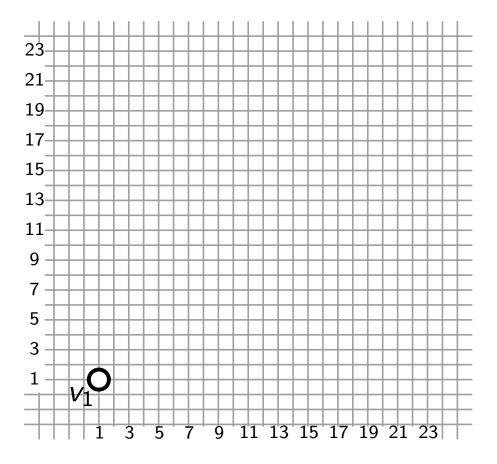


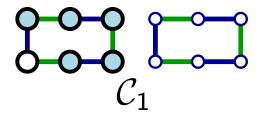


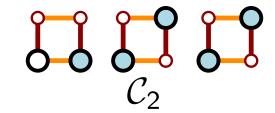


Placement algorithm:

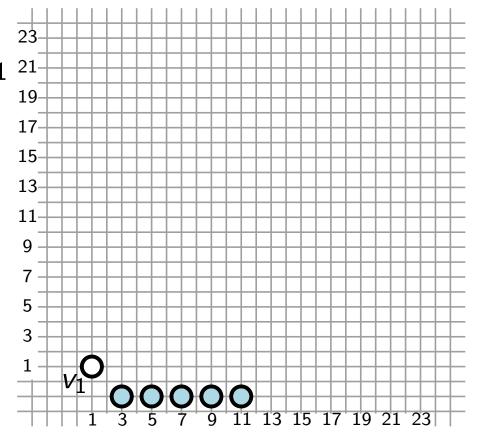
 $\bigcirc$  Pick  $v_1$ , place it

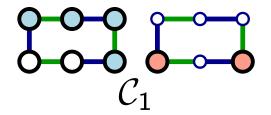


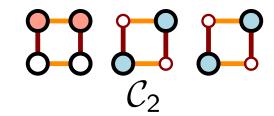




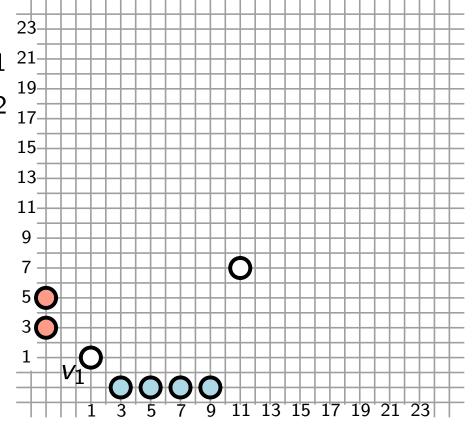
- $\bigcirc$  Pick  $v_1$ , place it
- lacktriangle Assign *x*-coords. to cycle in  $\mathcal{C}_1$

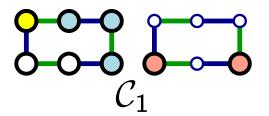


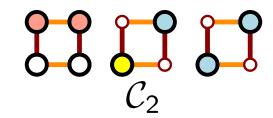




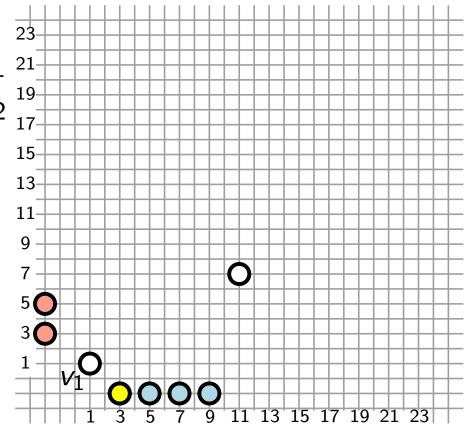
- $\bigcirc$  Pick  $v_1$ , place it
- lacktriangle Assign x-coords. to cycle in  $\mathcal{C}_1$
- lacktriangle Assign y-coords. to cycle in  $\mathcal{C}_2$

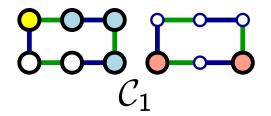


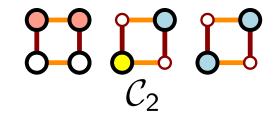




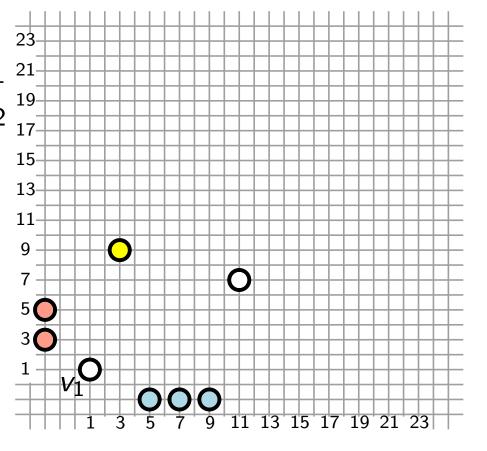
- $\bigcirc$  Pick  $v_1$ , place it
- Assign x-coords. to cycle in  $C_1$
- lacktriangle Assign y-coords. to cycle in  $\mathcal{C}_2$
- Pick v with 1 assigned coord.

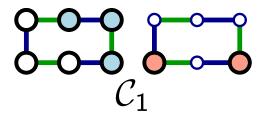


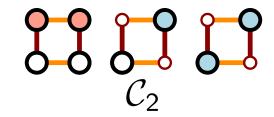




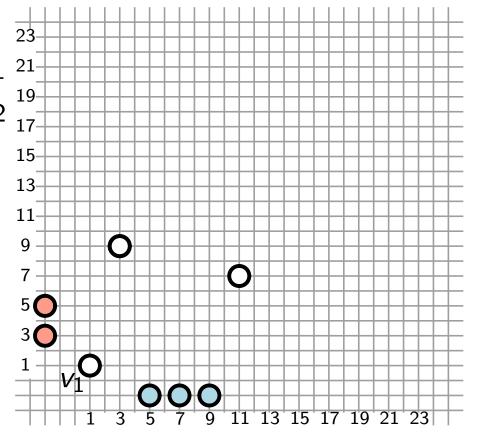
- $\bigcirc$  Pick  $v_1$ , place it
- Assign x-coords. to cycle in  $\mathcal{C}_1$
- lacktriangle Assign y-coords. to cycle in  $\mathcal{C}_2$
- Pick v with 1 assigned coord.
- Place v

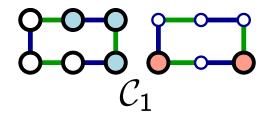


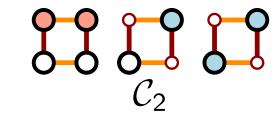




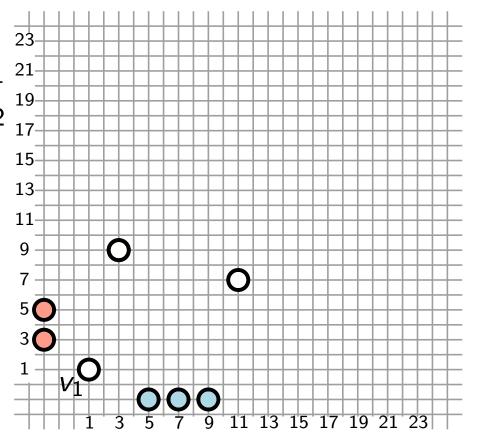
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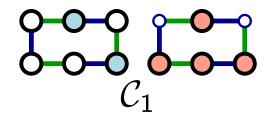


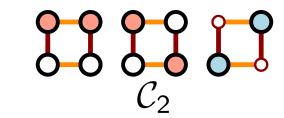




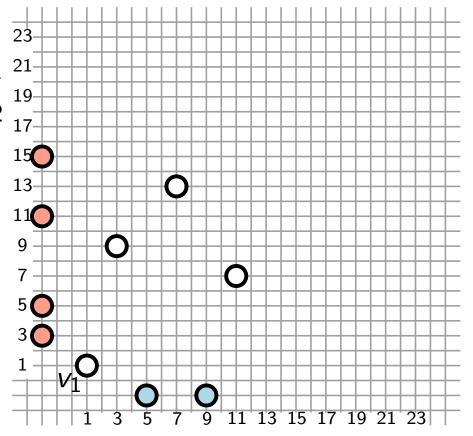
- ho Pick  $v_1$ , place it
- lack Assign x-coords. to cycle in  $\mathcal{C}_1$
- ullet Assign y-coords. to cycle in  $\mathcal{C}_2$
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- Place v

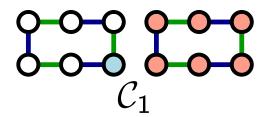


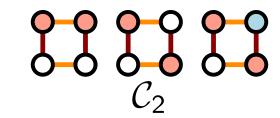




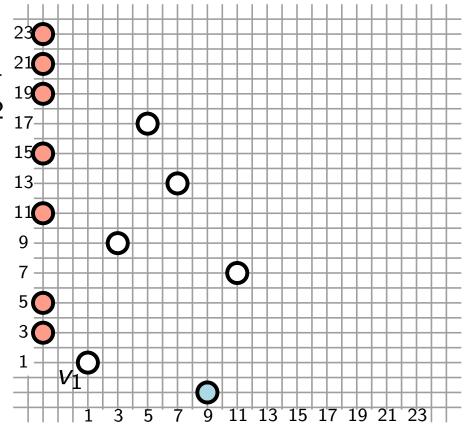
- m 
  ho Pick  $v_1$ , place it
- lack Assign x-coords. to cycle in  $\mathcal{C}_1$
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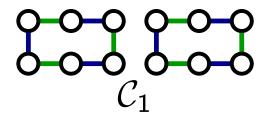


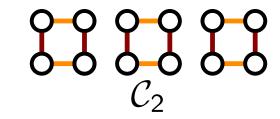




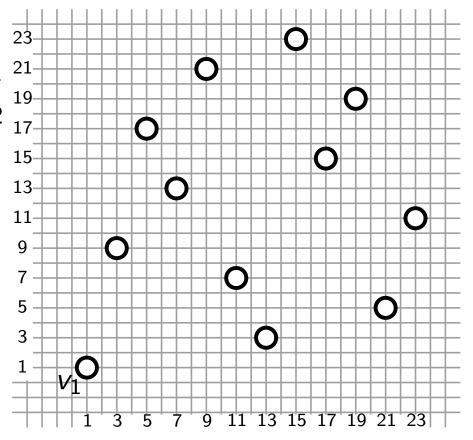
- ho Pick  $v_1$ , place it
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- Assign y-coords. to cycle in  $C_2$
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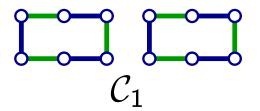


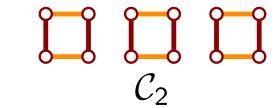




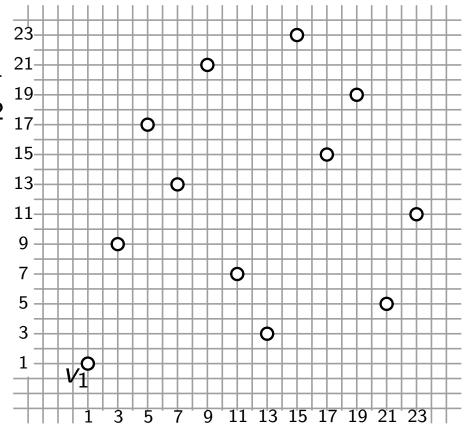
- ho Pick  $v_1$ , place it
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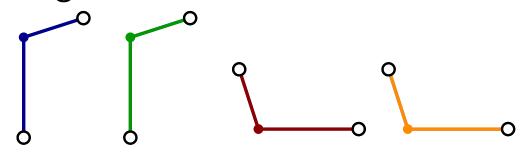


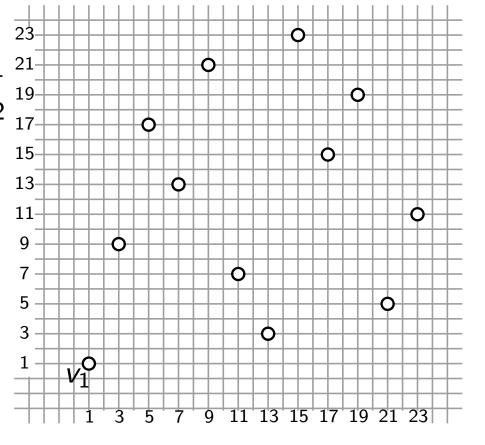


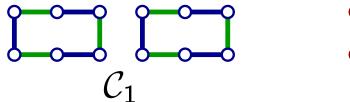
#### Placement algorithm:

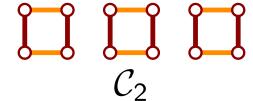
- ightharpoonup Pick  $v_1$ , place it
- Assign x-coords. to cycle in  $\mathcal{C}_1$
- ullet Assign y-coords. to cycle in  $\mathcal{C}_2$
- Pick v with 1 assigned coord.
- Place v

#### Edges:





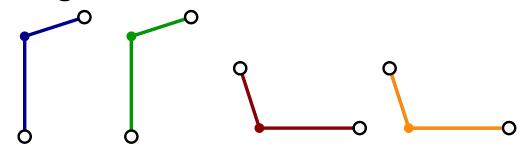


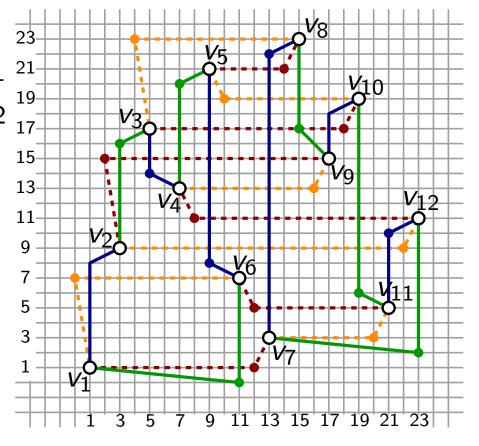


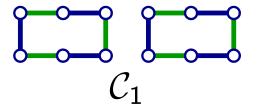
#### Placement algorithm:

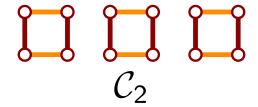
- ho Pick  $v_1$ , place it
- lack Assign x-coords. to cycle in  $\mathcal{C}_1$
- ullet Assign *y*-coords. to cycle in  $\mathcal{C}_2$
- Pick v with 1 assigned coord.
- Place v

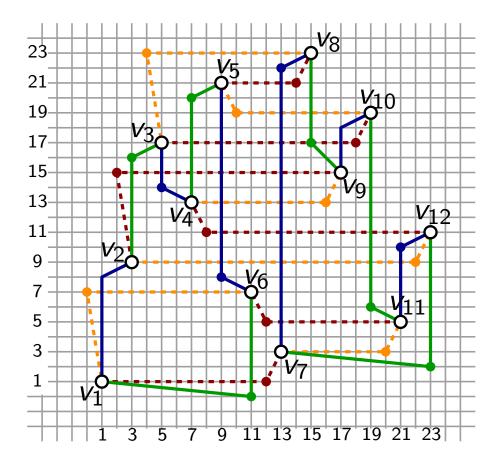
#### Edges:

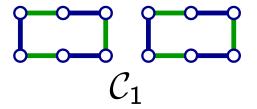


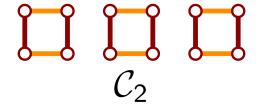


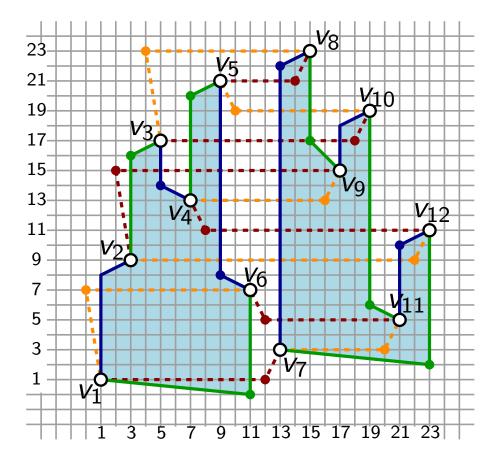


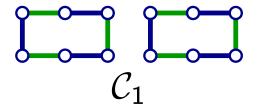


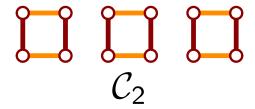


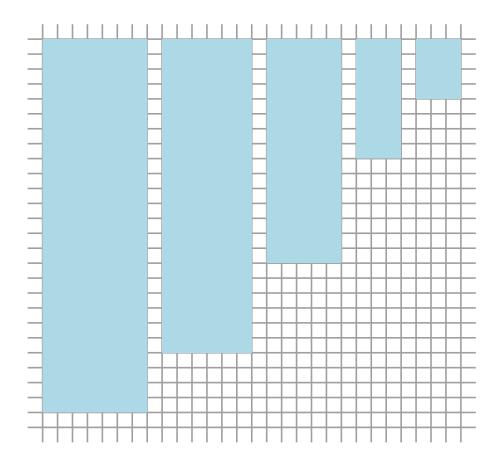


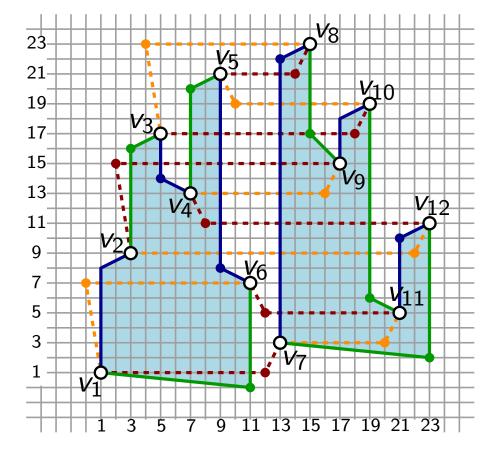


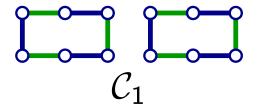


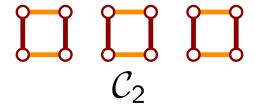


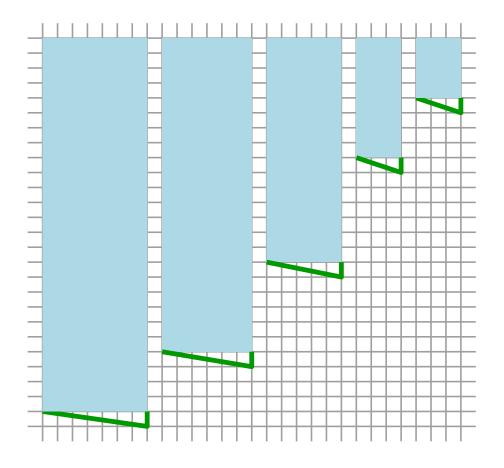


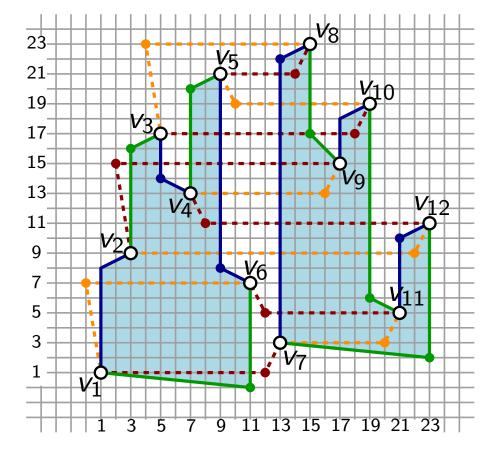


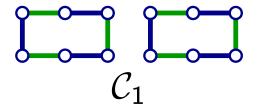


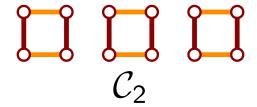


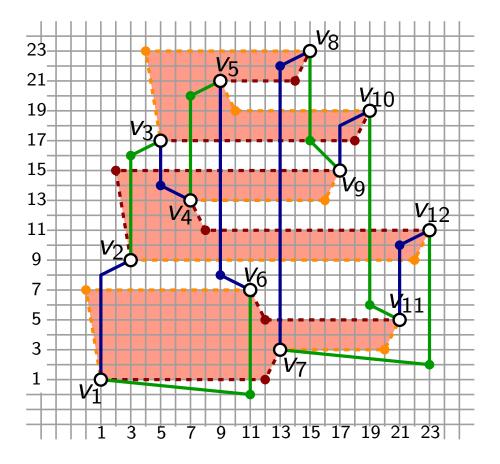


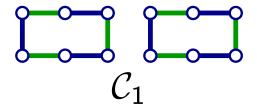


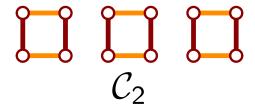


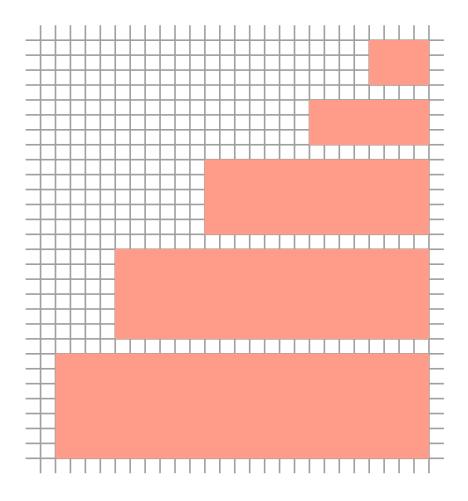


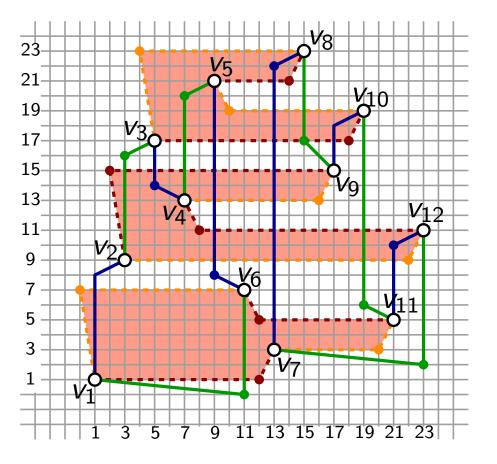


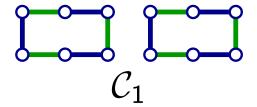


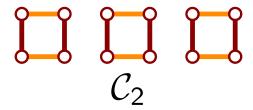


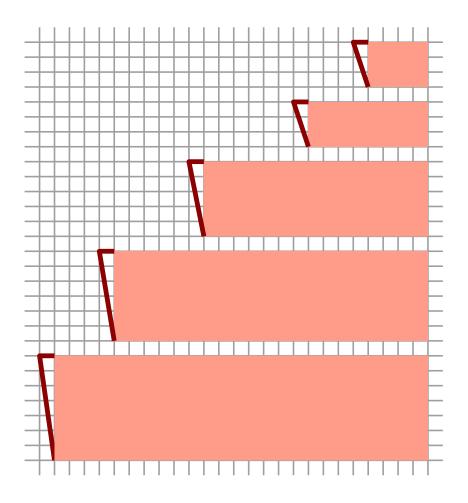


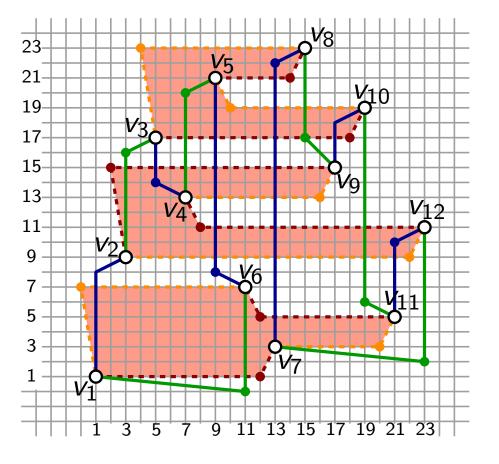


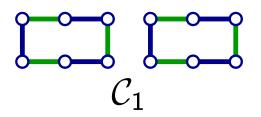


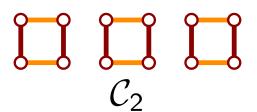












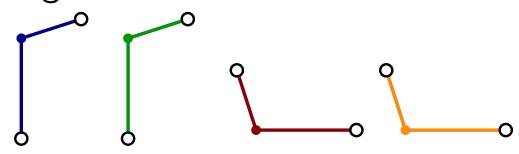
Bends:  $1 \times 1$ 

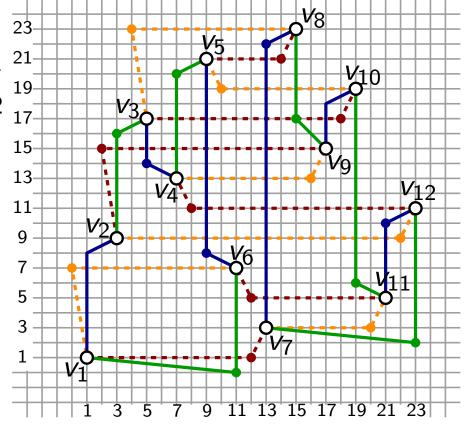
Grid size:  $2n \times 2n$ 

#### Placement algorithm:

- ho Pick  $v_1$  , place it
  - Assign *x*-coords. to cycle in  $\mathcal{C}_1$
- lacktriangle Assign *y*-coords. to cycle in  $\mathcal{C}_2$
- Pick v with 1 assigned coord.
- Place v

#### Edges:





## Overview

Graph classes		Number of bends	
Cycle	×	Cycle	1  imes 1
Caterpillar	×	Cycle	1  imes 1
Four Matchings			1  imes 1  imes 1  imes 1
Tree	X	Matching	$1 \times 0$
Wheel	X	Matching	$2 \times 0$
Outerpath	×	Matching	$2 \times 1$
Outerplanar	X	Outerplanar	$3 \times 3$
2-page book emb.	×	2-page book emb.	$4 \times 4$
Planar	×	Planar	$6 \times 6$

Idea: 

Matching edges horizontal, tree edges with 1 bend

Idea: O Matching edges horizontal, tree edges with 1 bend

• Place matching edges inductively  $\Rightarrow$  *y*-coord.

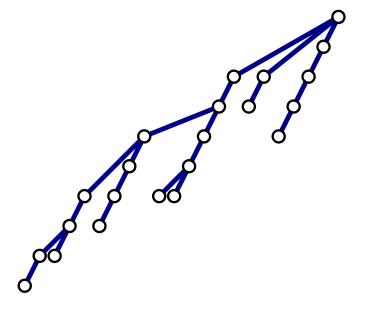
Idea: 
Matching edges horizontal, tree edges with 1 bend

• Place matching edges inductively  $\Rightarrow$  *y*-coord.

• Use post-order on tree  $\Rightarrow$  *x*-coord.

Idea: O Matching edges horizontal, tree edges with 1 bend

- Place matching edges inductively  $\Rightarrow$  *y*-coord.
- Use post-order on tree  $\Rightarrow x$ -coord.



Idea: O Matching edges horizontal, tree edges with 1 bend

• Place matching edges inductively  $\Rightarrow$  *y*-coord.

• Use post-order on tree  $\Rightarrow x$ -coord.

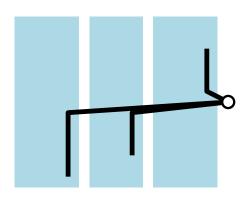
 $\circ$   $\Rightarrow$  Subtrees in disjoint *x*-intervals

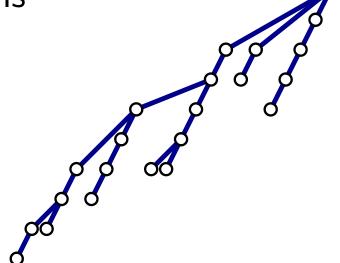
Idea: O Matching edges horizontal, tree edges with 1 bend

• Place matching edges inductively  $\Rightarrow$  *y*-coord.

• Use post-order on tree  $\Rightarrow x$ -coord.

 $\circ$   $\Rightarrow$  Subtrees in disjoint *x*-intervals





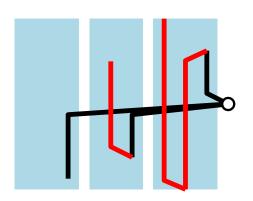
Idea: O Matching edges horizontal, tree edges with 1 bend

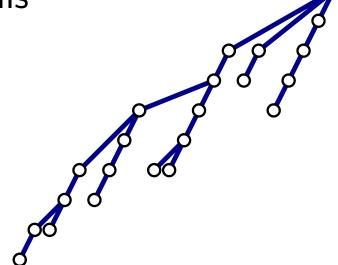
• Place matching edges inductively  $\Rightarrow$  *y*-coord.

• Use post-order on tree  $\Rightarrow x$ -coord.

 $\circ$   $\Rightarrow$  Subtrees in disjoint *x*-intervals

#### Problem:





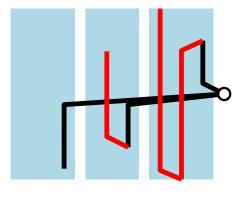
Idea: O Matching edges horizontal, tree edges with 1 bend

 $\bigcirc$  Place matching edges inductively  $\Rightarrow y$ -coord.

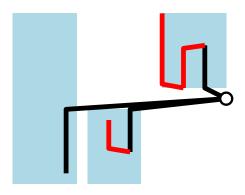
• Use post-order on tree  $\Rightarrow x$ -coord.

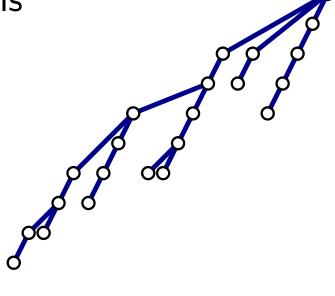
 $\circ$   $\Rightarrow$  Subtrees in disjoint *x*-intervals

Problem:



Solution:





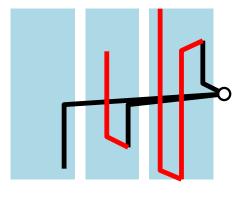
Idea: O Matching edges horizontal, tree edges with 1 bend

ullet Place matching edges inductively  $\Rightarrow y$ -coord.

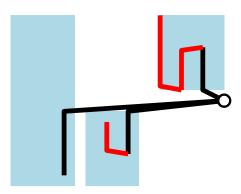
• Use post-order on tree  $\Rightarrow x$ -coord.

 $\circ$   $\Rightarrow$  Subtrees in disjoint *x*-intervals

Problem:



Solution:



All but one subtree completely above or completely below.

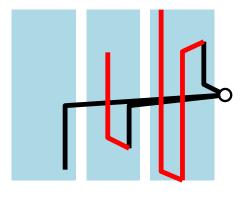
Idea: O Matching edges horizontal, tree edges with 1 bend

• Place matching edges inductively  $\Rightarrow$  *y*-coord.

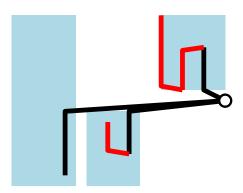
• Use post-order on tree  $\Rightarrow x$ -coord.

 $\circ$   $\Rightarrow$  Subtrees in disjoint *x*-intervals

Problem:



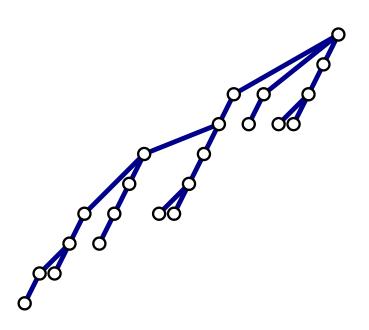
Solution:

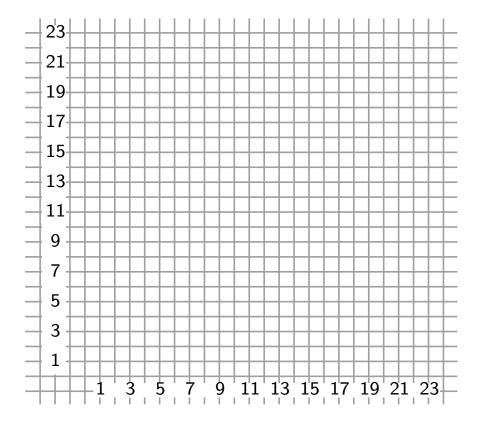


All but one subtree completely above or completely below.

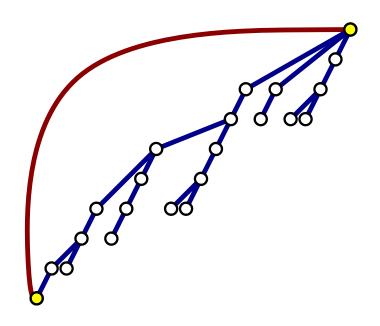
Main ideas adopted from [Cabello et al. JGAA'11, Di Giacomo et al. JGAA'09].

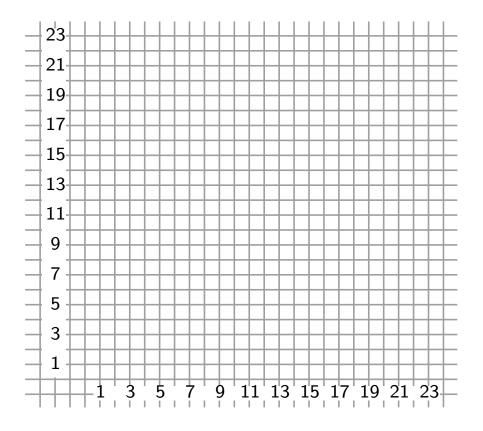
Place root + matching at the top



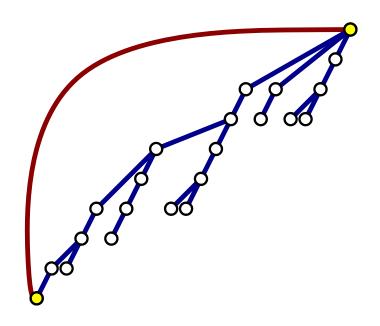


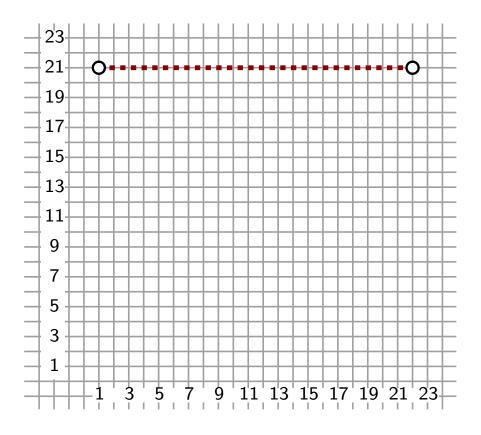
Place root + matching at the top



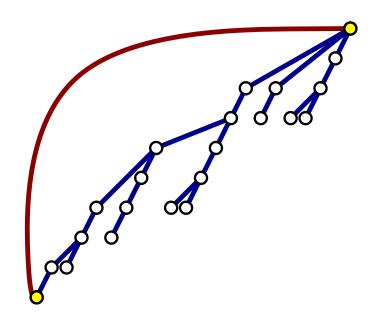


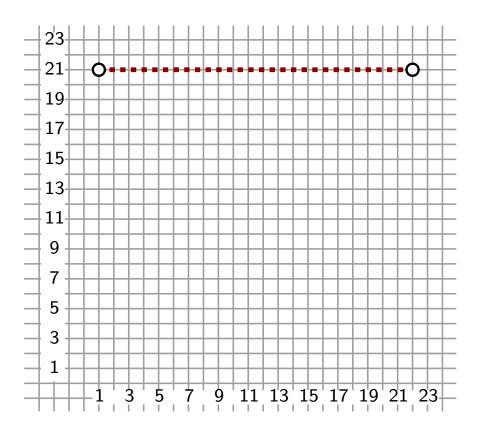
Place root + matching at the top



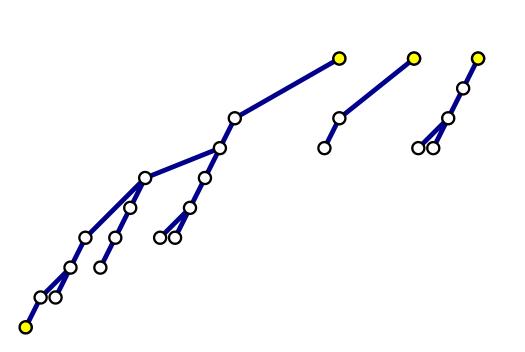


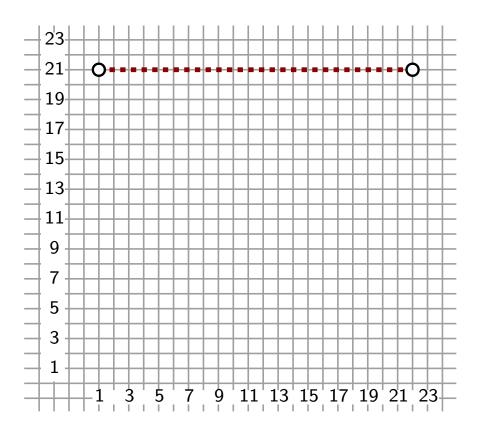
- Place root + matching at the top
- Split the tree



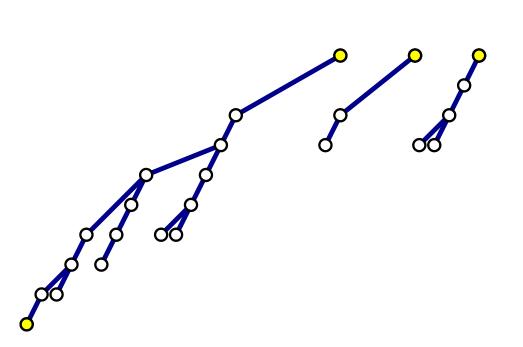


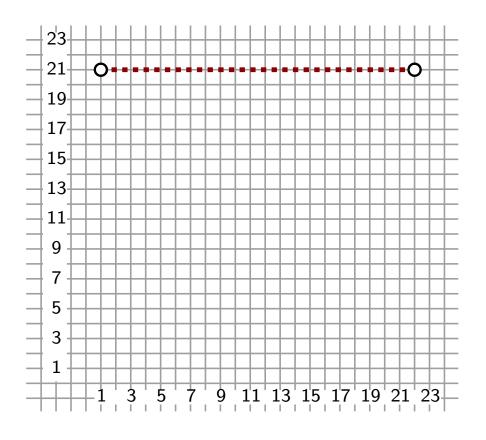
- Place root + matching at the top
- Split the tree



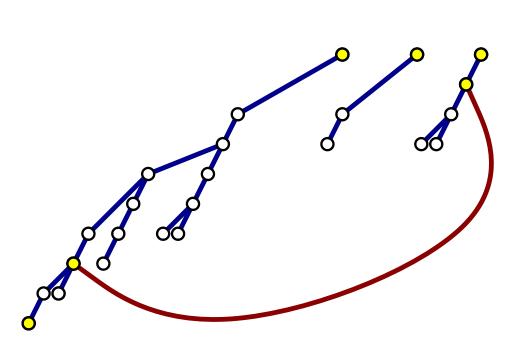


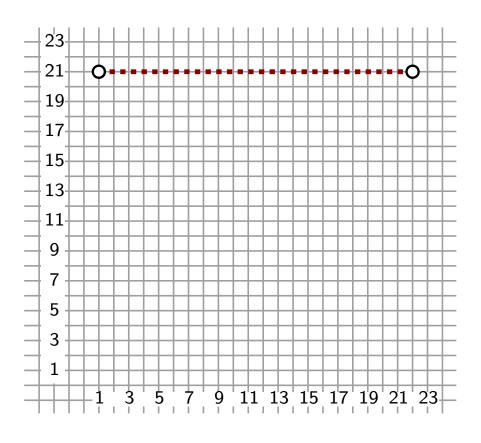
- Place root + matching at the top
- Split the tree
- Place vertex adj. to placed vertex (+ matching) at the top



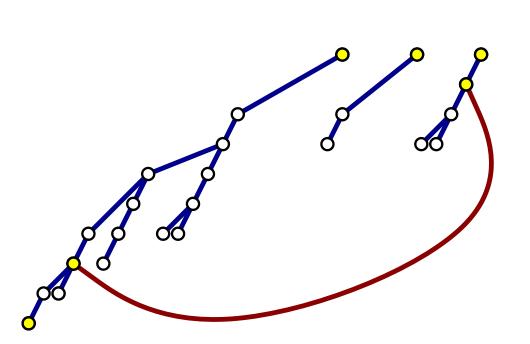


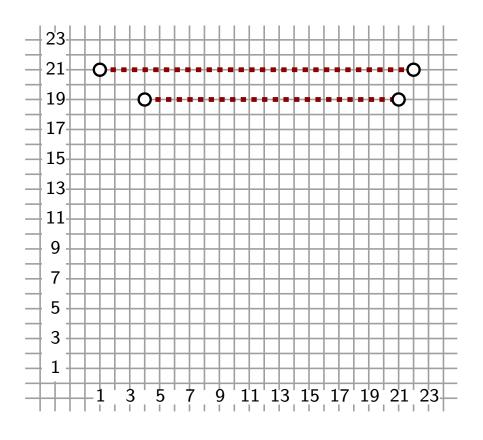
- Place root + matching at the top
- Split the tree
- Place vertex adj. to placed vertex (+ matching) at the top



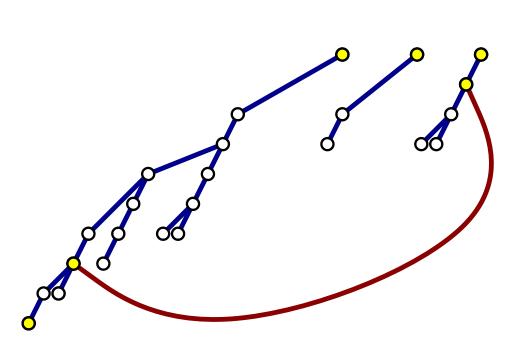


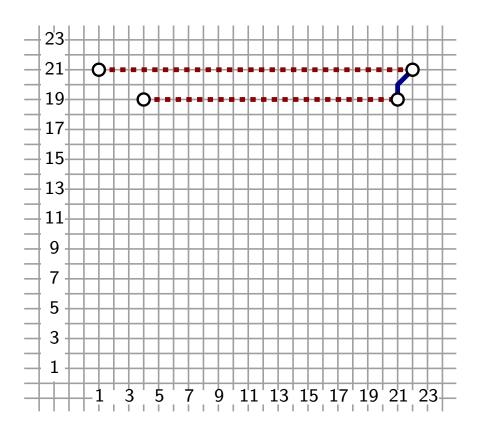
- Place root + matching at the top
- Split the tree
- Place vertex adj. to placed vertex (+ matching) at the top



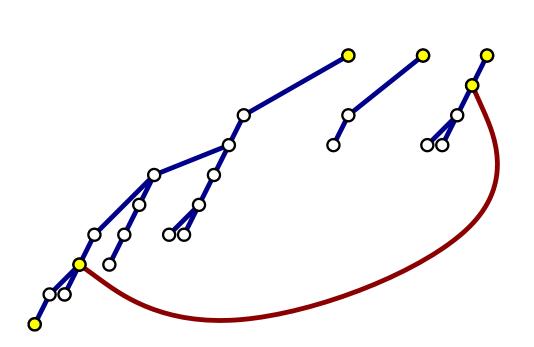


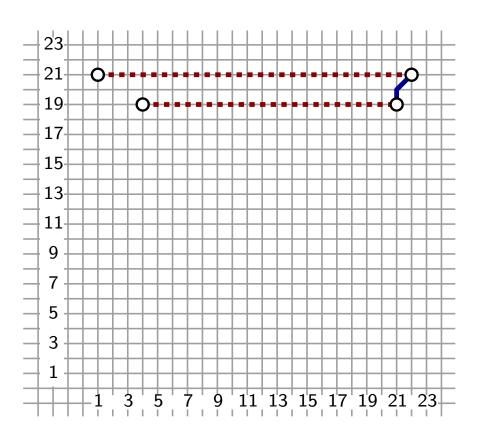
- Place root + matching at the top
- Split the tree
- Place vertex adj. to placed vertex (+ matching) at the top



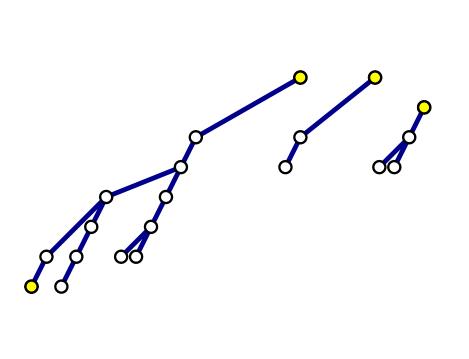


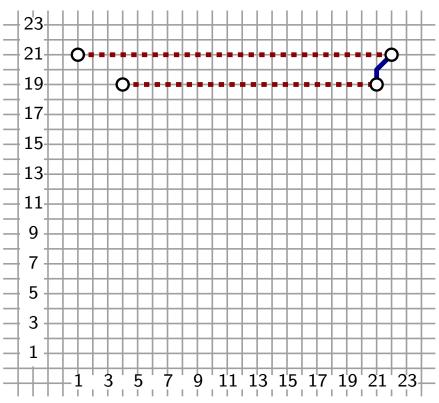
- Place root + matching at the top
- ►○ Split the tree
- - $\circ$  Place vertex adj. to placed vertex (+ matching) at the top





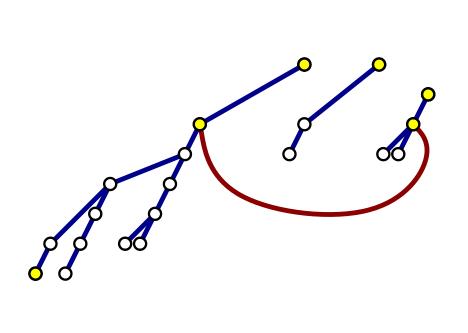
- Place root + matching at the top
- ►○ Split the tree
- - $\circ$  Place vertex adj. to placed vertex (+ matching) at the top

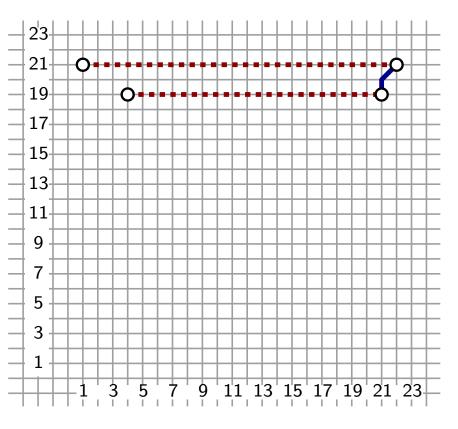




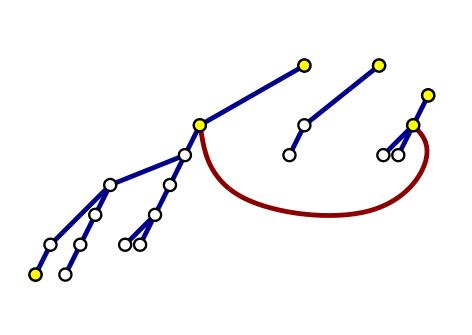


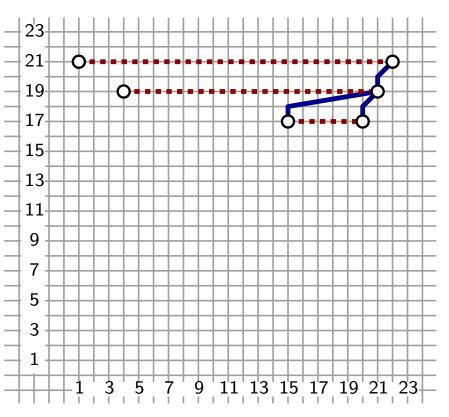
- Place root + matching at the top
- ►○ Split the tree
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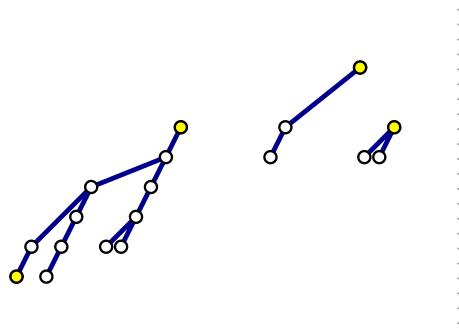


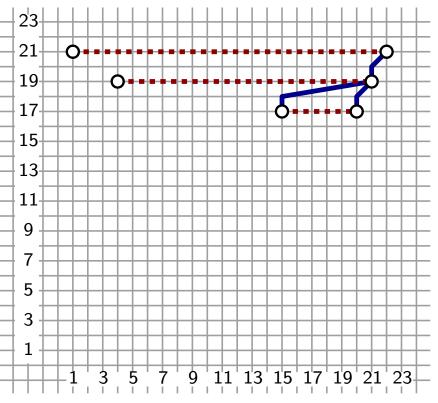
- Place root + matching at the top
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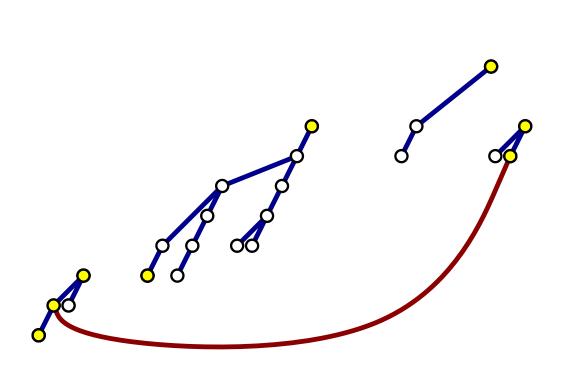


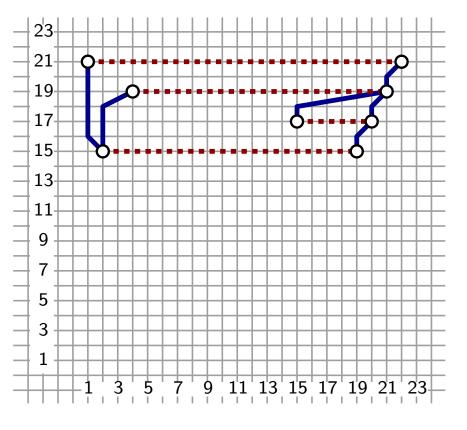
- Place root + matching at the top
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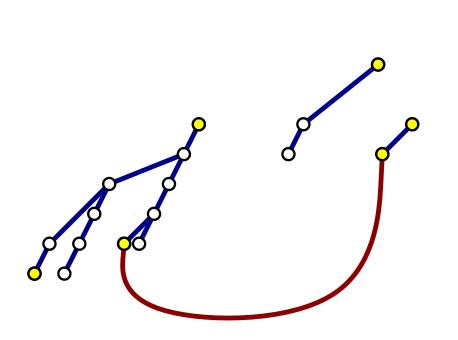


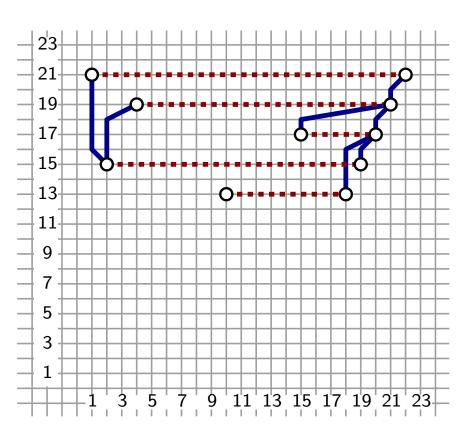
- Place root + matching at the top
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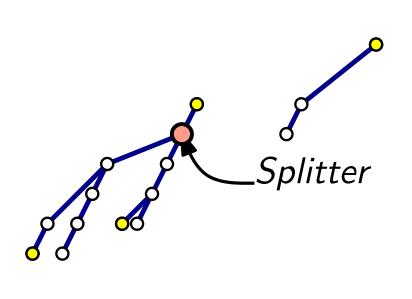


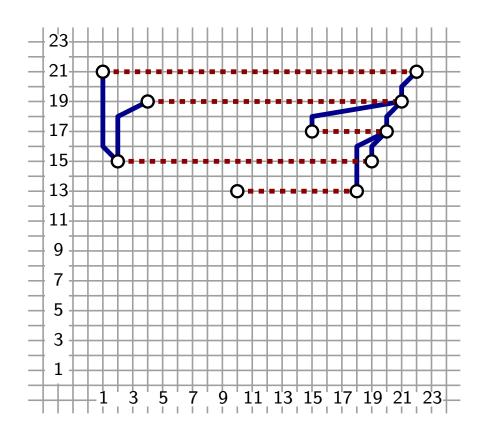
- Place root + matching at the top
- ►○ Split the tree
- - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top



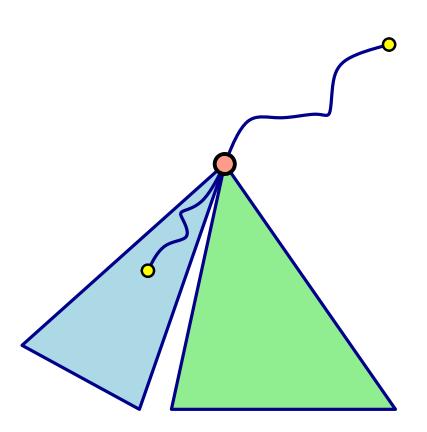


- Place root + matching at the top
- Split the tree
  - - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top

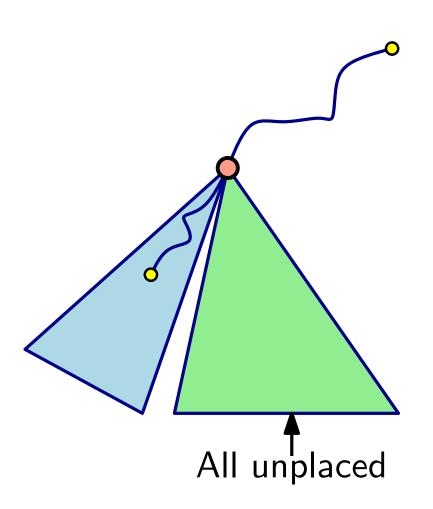




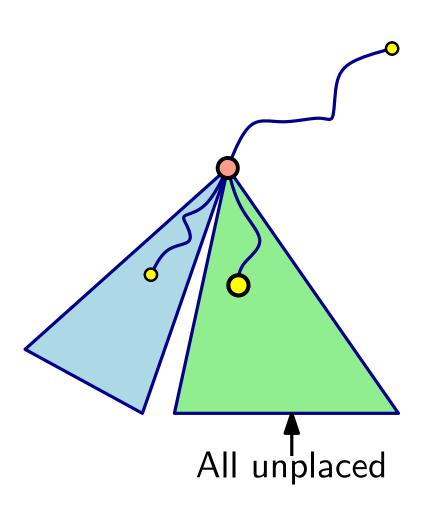
- Place root + matching at the top
- Split the tree
  - Place vertex adj. to placed vertex (+ matching) at the top



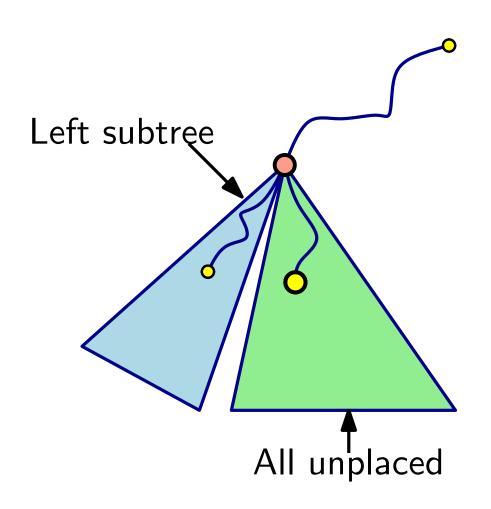
- Place root + matching at the top
- Split the tree
  - - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top



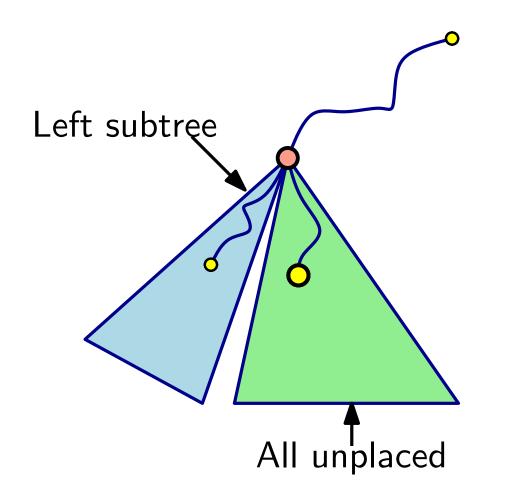
- Place root + matching at the top
- Split the tree
  - - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top



- Place root + matching at the top
- Split the tree
  - - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top



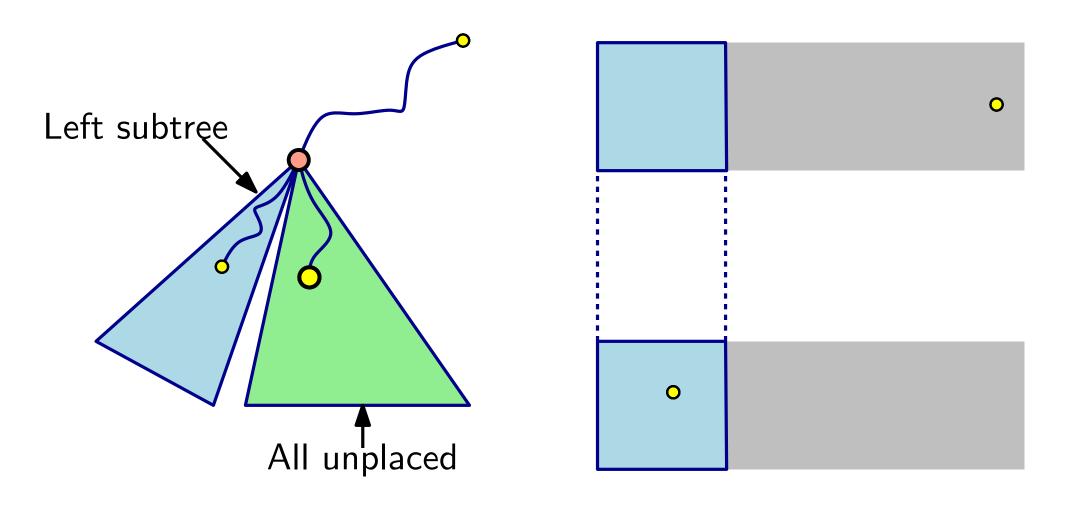
- Place root + matching at the top
- Split the tree
  - - $\circ$  Place vertex adj. to placed vertex (+ matching) at the top



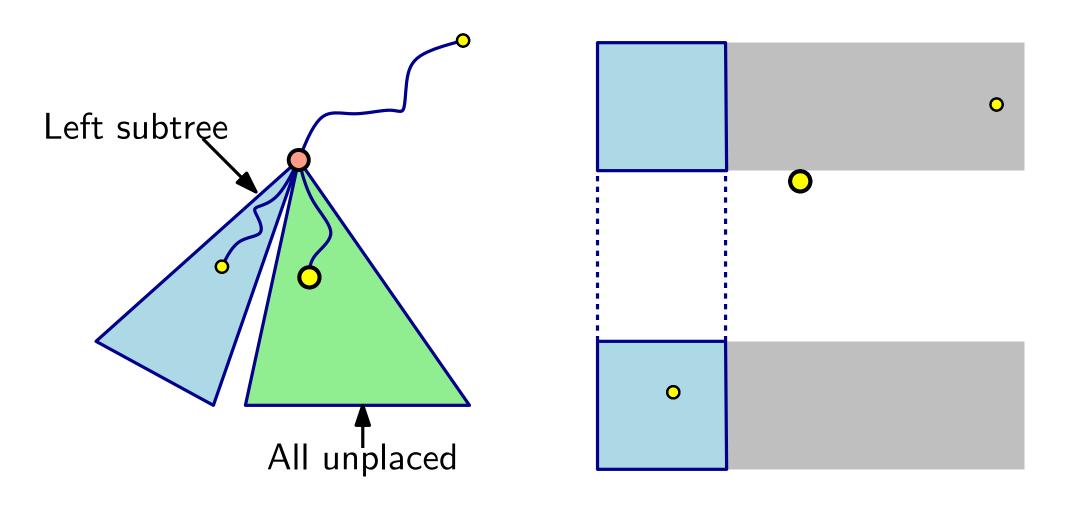
Placed vertices

Placed vertices

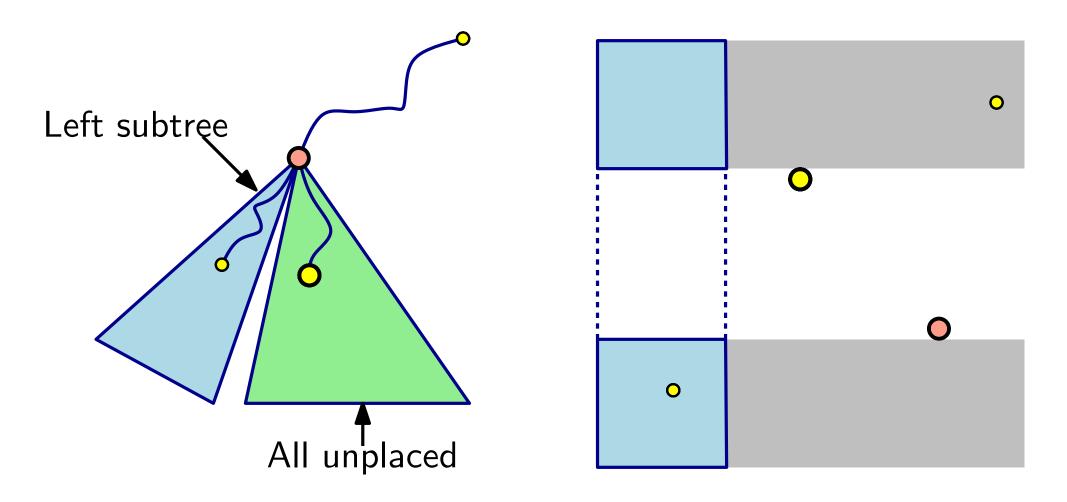
- Place root + matching at the top
- Split the tree
  - - $\circ$  Place vertex adj. to placed vertex (+ matching) at the top



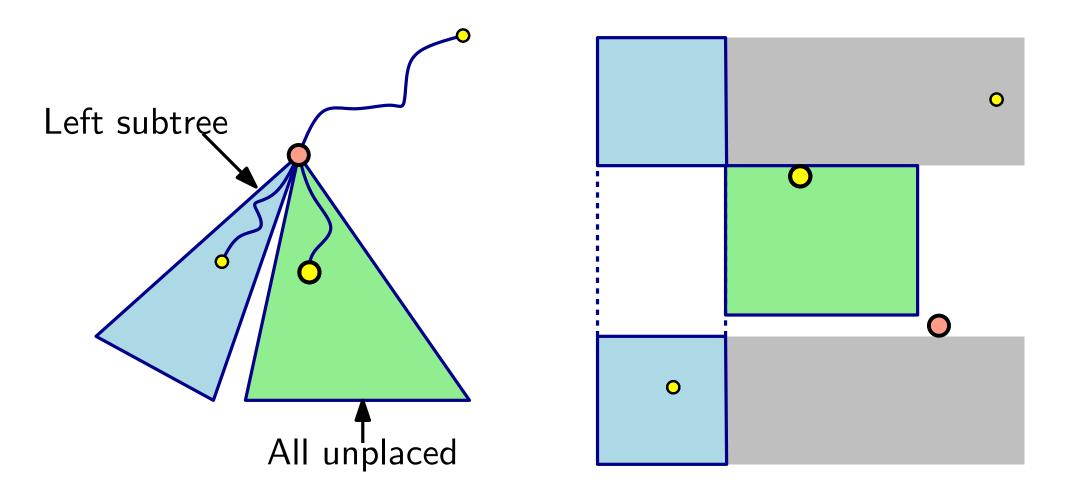
- Place root + matching at the top
- Split the tree
  - - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top



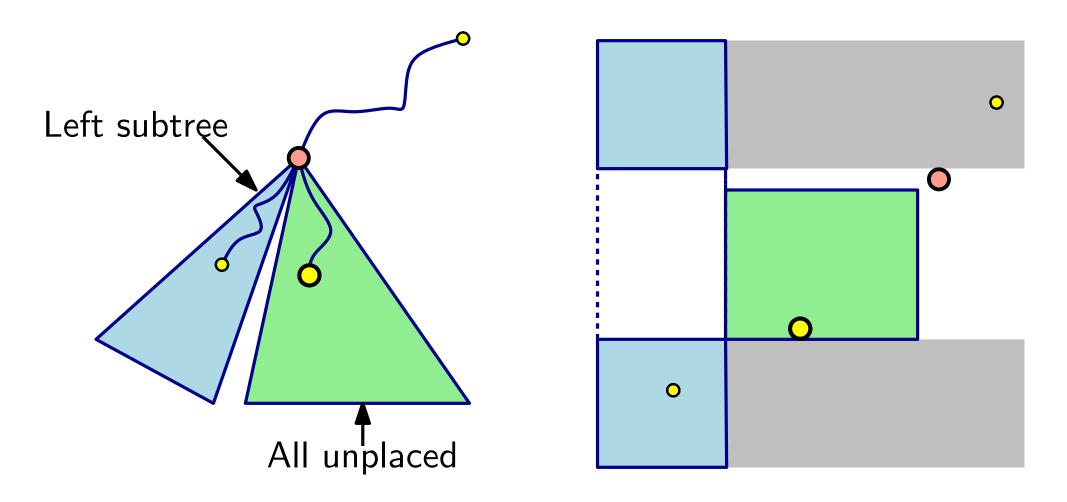
- Place root + matching at the top
- Split the tree
  - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top



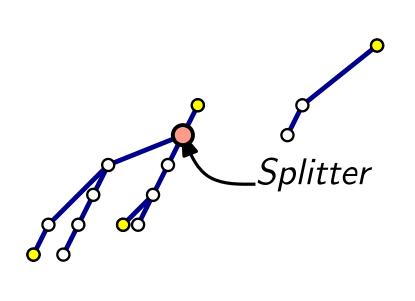
- Place root + matching at the top
- Split the tree
  - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top

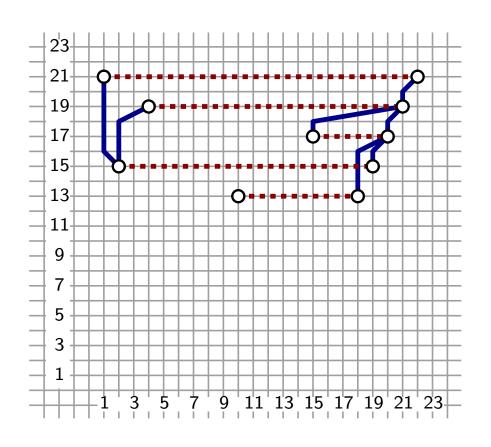


- Place root + matching at the top
- Split the tree
  - - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top

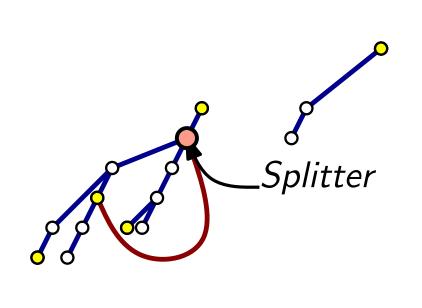


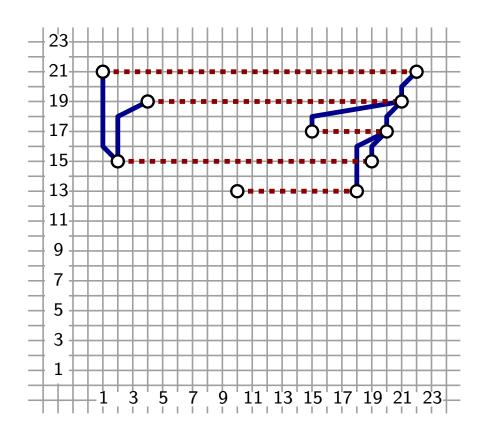
- Place root + matching at the top
- ►○ Split the tree
- $\cdot$  Place vertex adj. to placed vertex (+ matching) at the top
- If splitter: place on opposite side



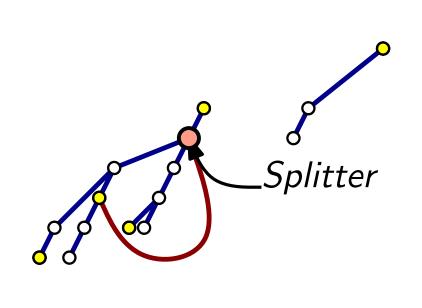


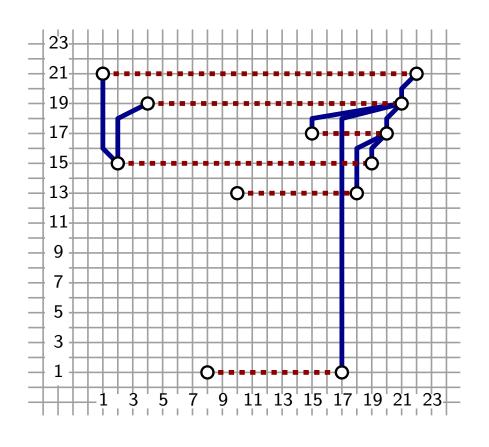
- Place root + matching at the top
- ►○ Split the tree
- - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top
- If splitter: place on opposite side



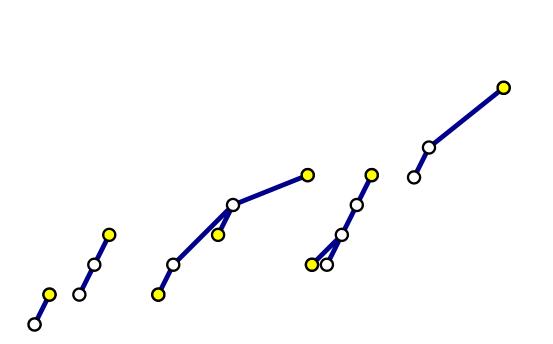


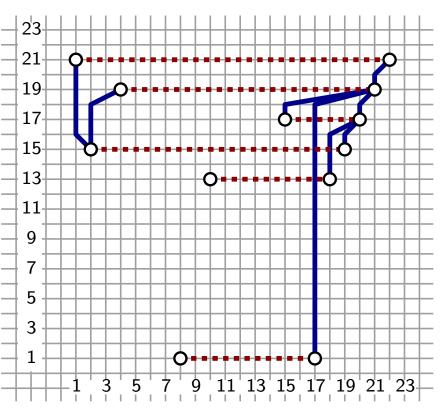
- Place root + matching at the top
- **➣** Split the tree
  - - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top
  - If splitter: place on opposite side



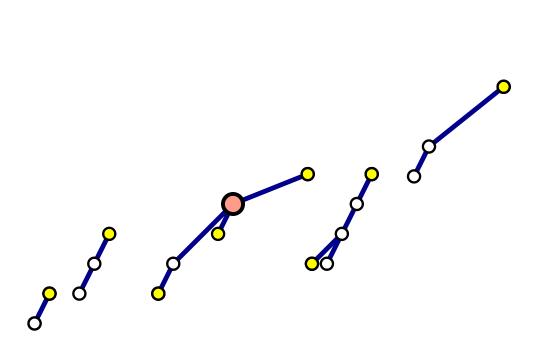


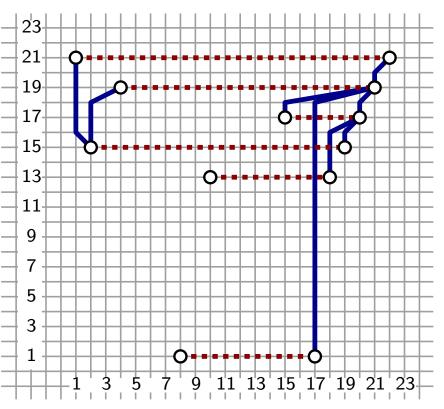
- Place root + matching at the top
- ►○ Split the tree
- $\cdot$  Place vertex adj. to placed vertex (+ matching) at the top
- If splitter: place on opposite side



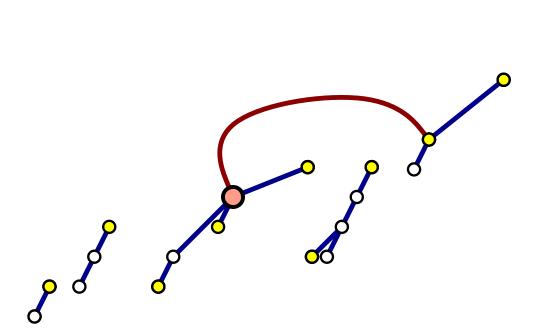


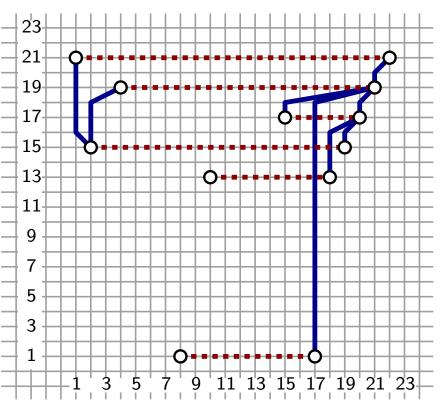
- Place root + matching at the top
- ►○ Split the tree
  - - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top
  - If splitter: place on opposite side



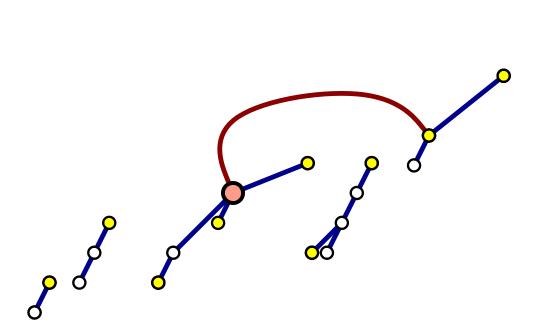


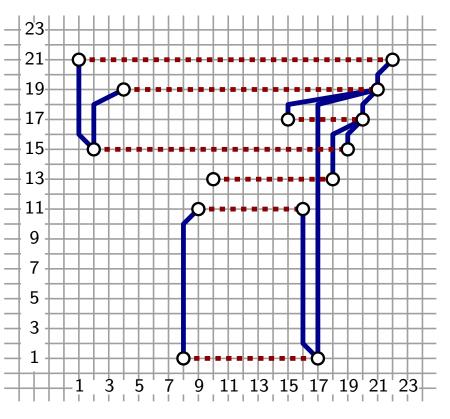
- Place root + matching at the top
- ►○ Split the tree
  - - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top
  - If splitter: place on opposite side



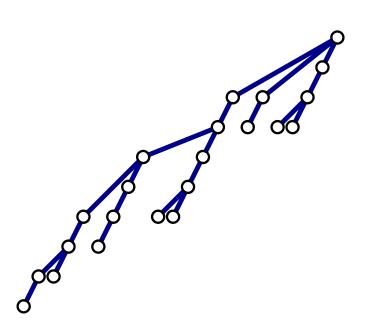


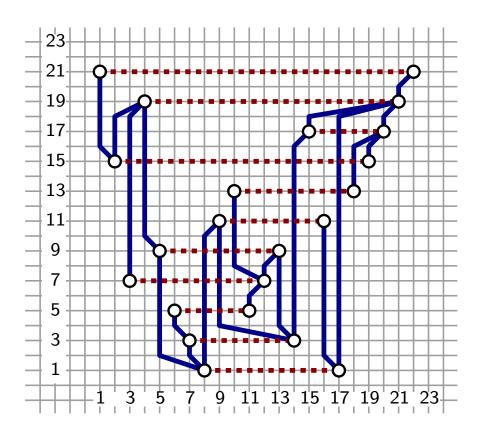
- Place root + matching at the top
- ►○ Split the tree
  - - $\bigcirc$  Place vertex adj. to placed vertex (+ matching) at the top
  - If splitter: place on opposite side





- Place root + matching at the top
- Split the tree
  Place vertex adj. to placed vertex (+ matching) at the top
  - If splitter: place on opposite side

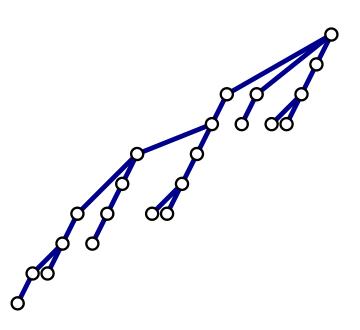


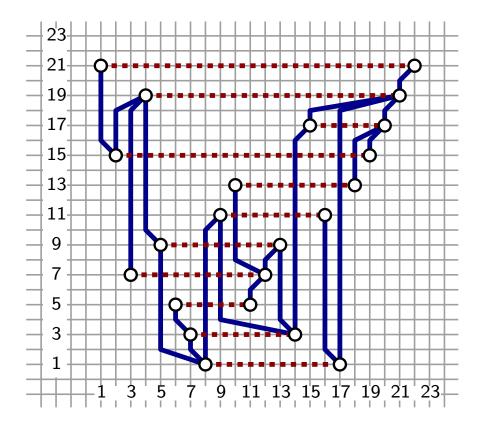


- Place root + matching at the top
- Split the tree
  Place vertex adj. to placed vertex (+ matching) at the top
  - If splitter: place on opposite side

Bends:  $1 \times 0$ 

Grid size:  $n \times (n-1)$ 



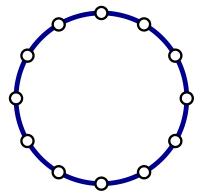


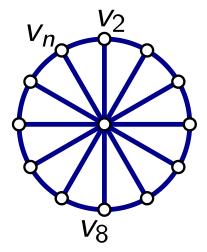
#### Overview

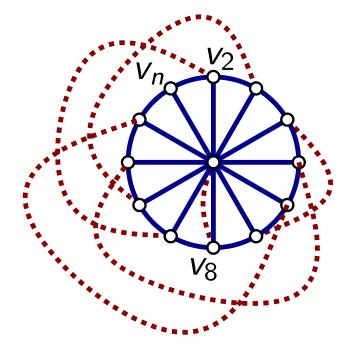
#### Graph classes

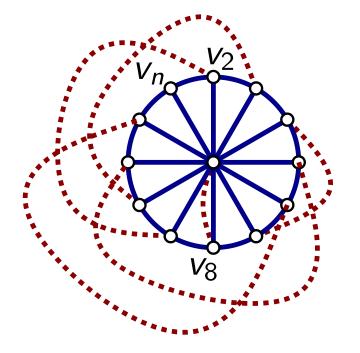
#### Number of bends

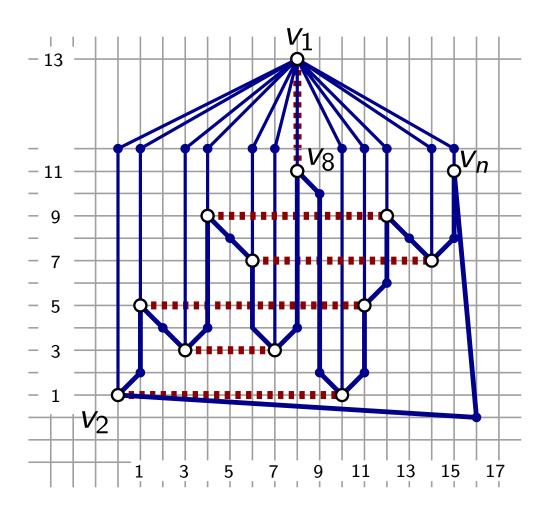
Cycle	X	Cycle	1  imes 1
Caterpillar	×	Cycle	1  imes 1
Four Matchings			1  imes 1  imes 1  imes 1
Tree	×	Matching	1  imes 0
Wheel	×	Matching	$2 \times 0$
Outerpath	×	Matching	2  imes 1
Outerplanar	×	Outerplanar	$3 \times 3$
2-page book emb.	×	2-page book emb.	$4 \times 4$
Planar	×	Planar	$6 \times 6$

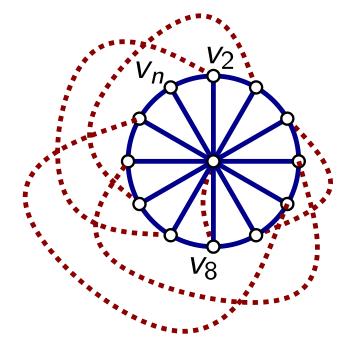






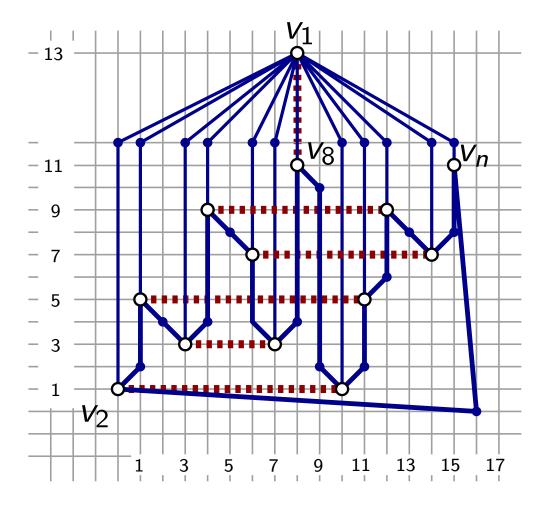


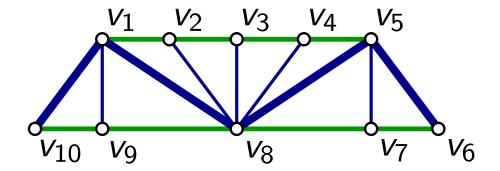




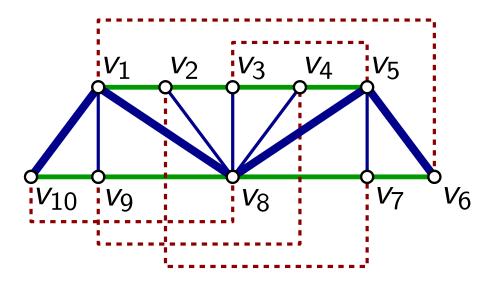
Bends:  $2 \times 0$ 

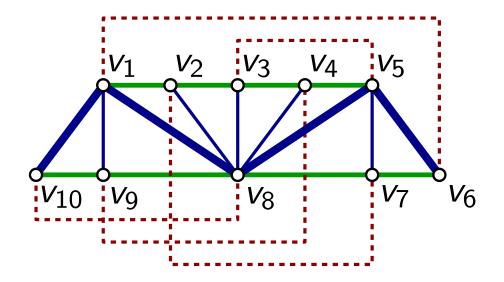
Grid size:  $(1.5n - 1) \times (n + 2)$ 

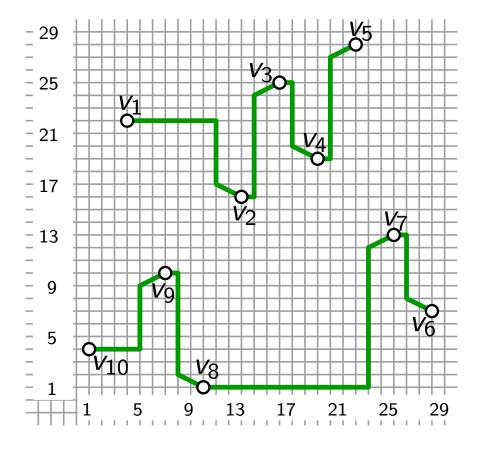


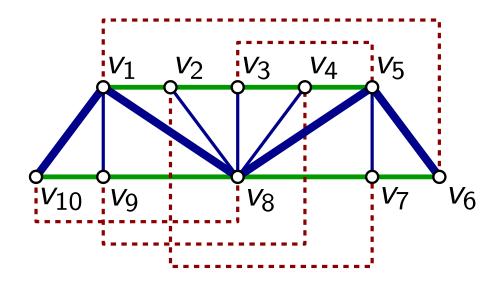


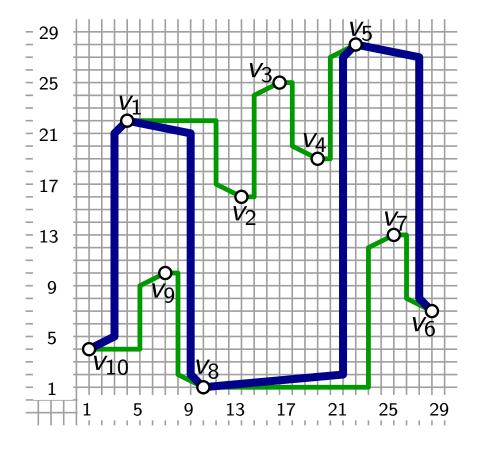
# Outerpath $\times$ Matching

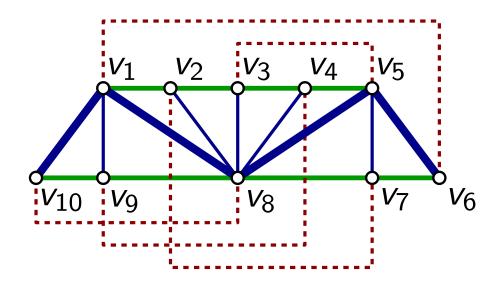


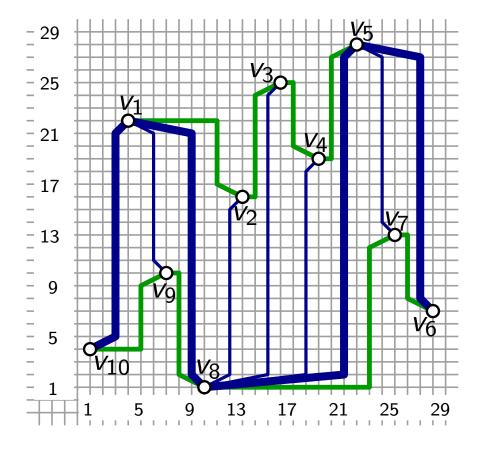


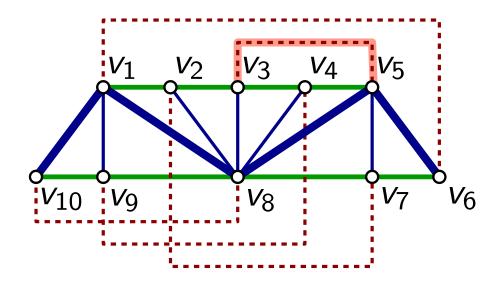


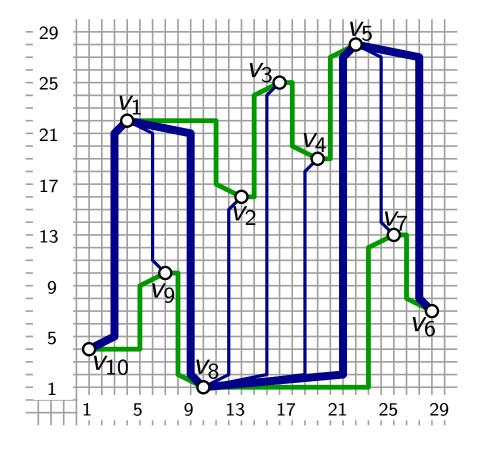


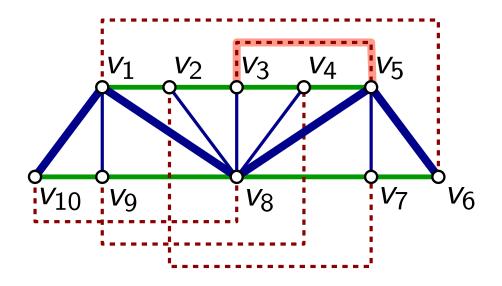


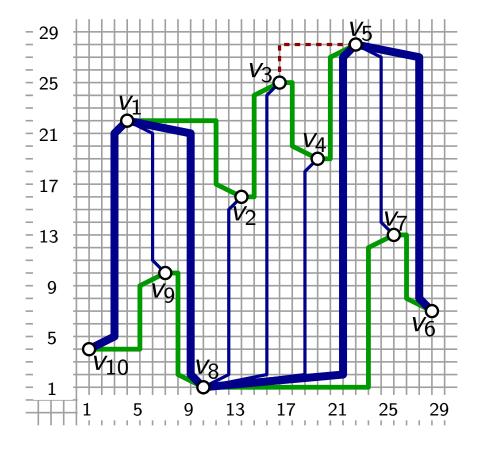


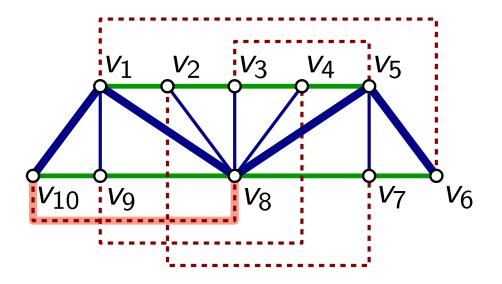


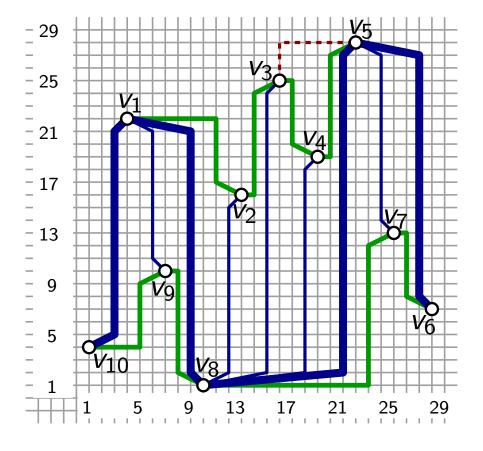


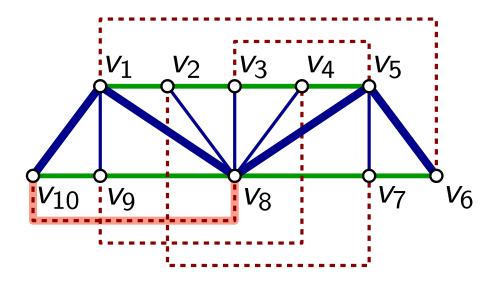


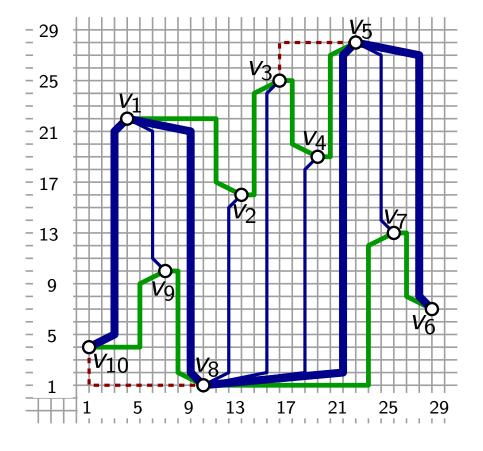


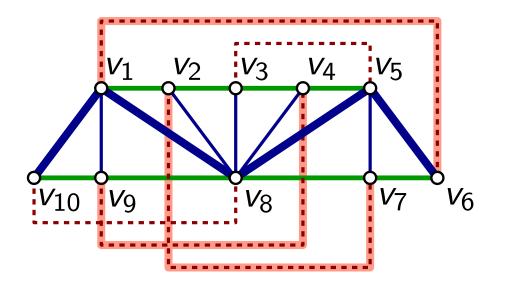


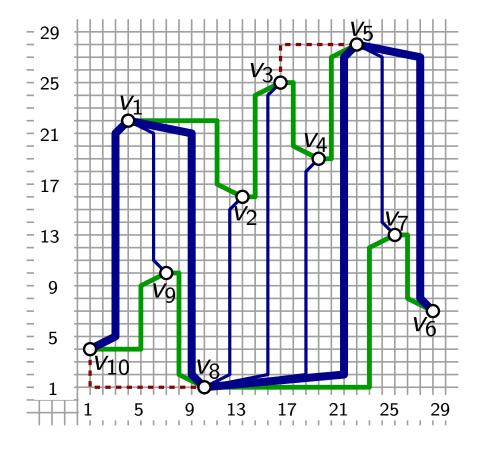


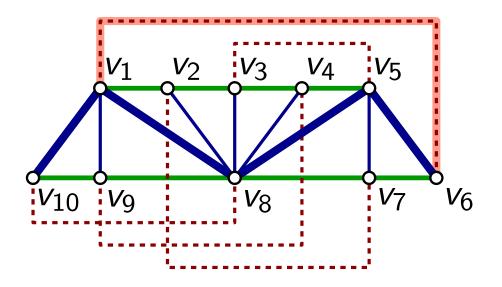


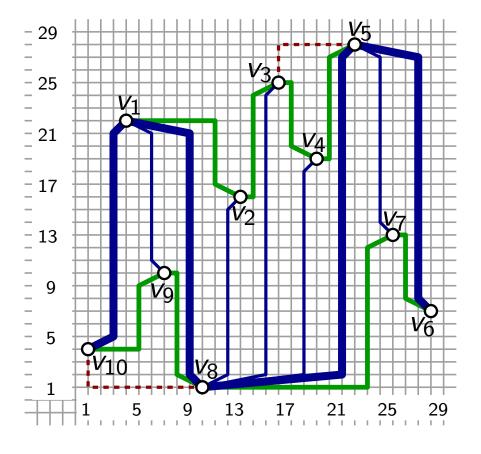


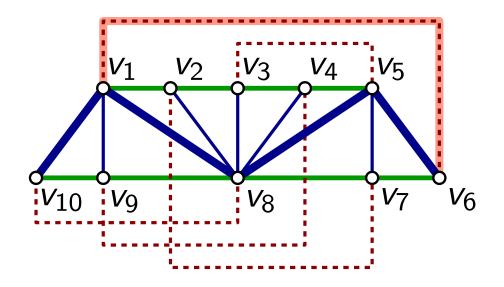


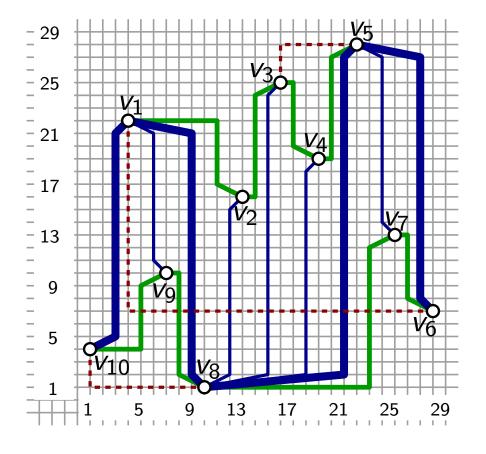


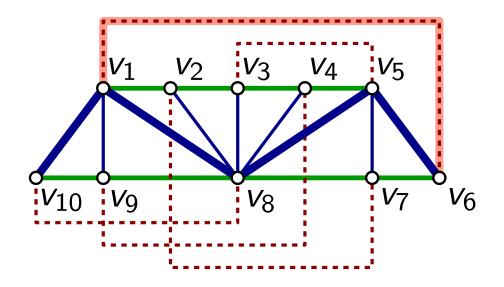


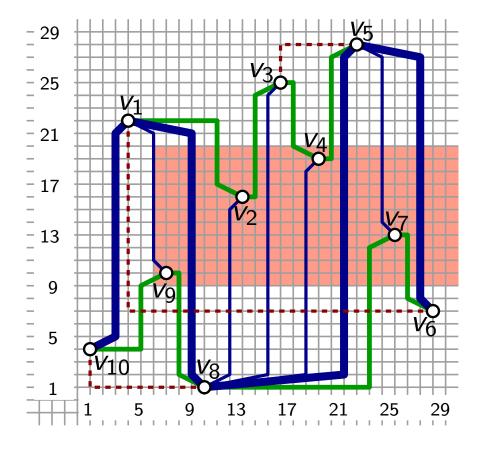


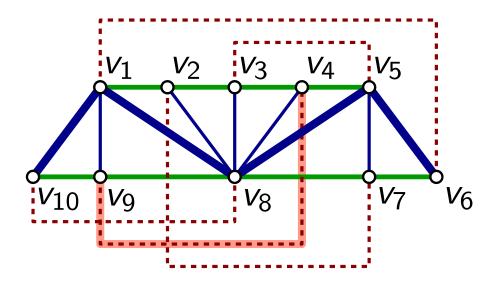


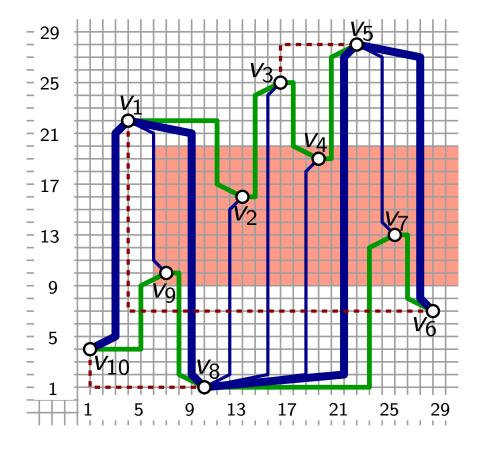


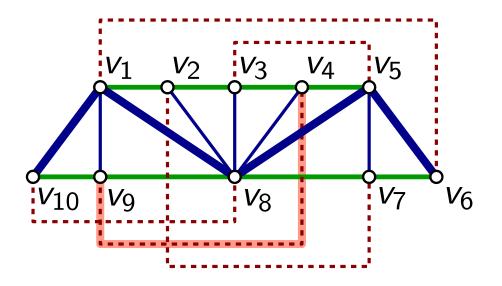


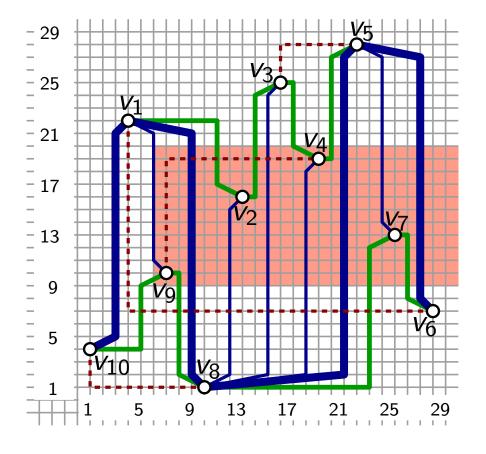


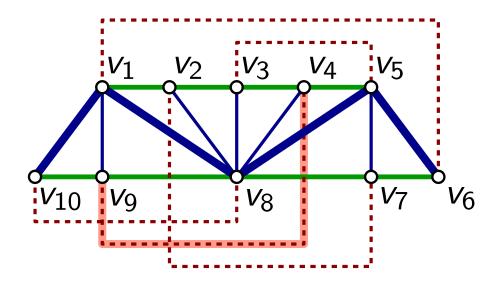


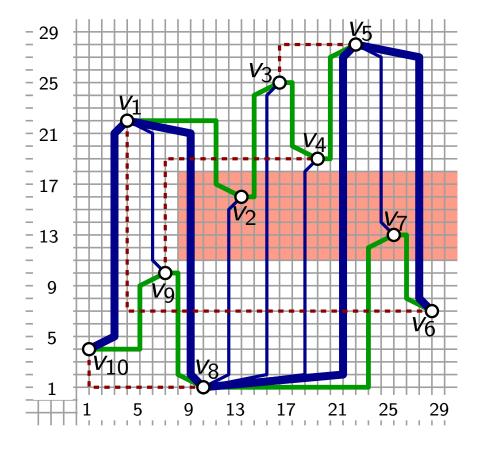


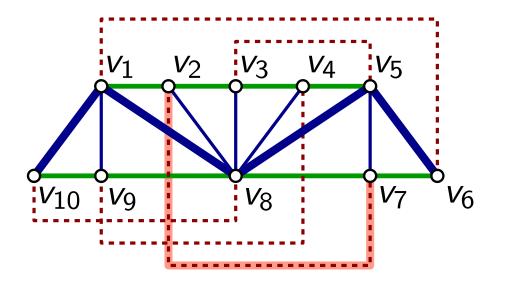


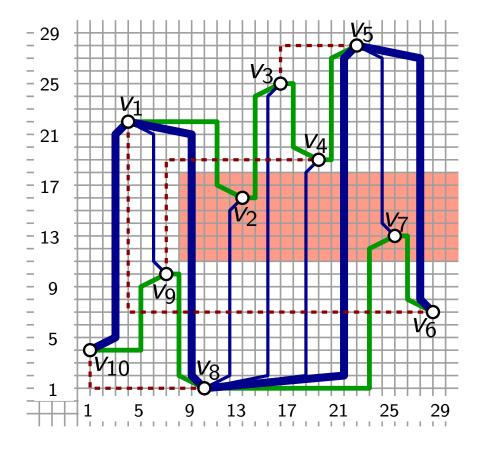


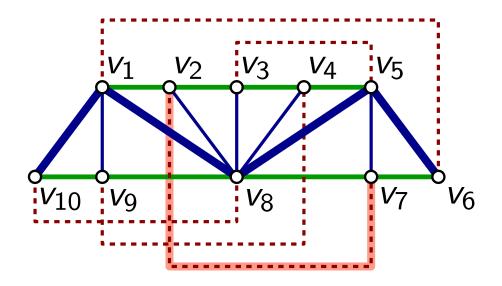


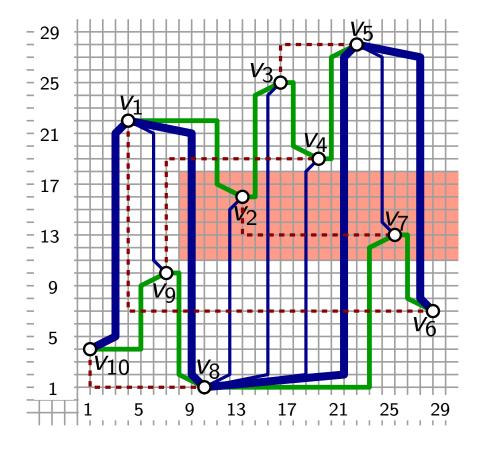


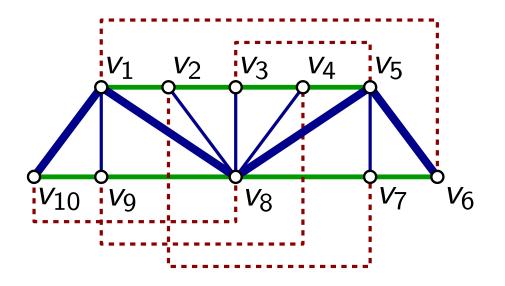


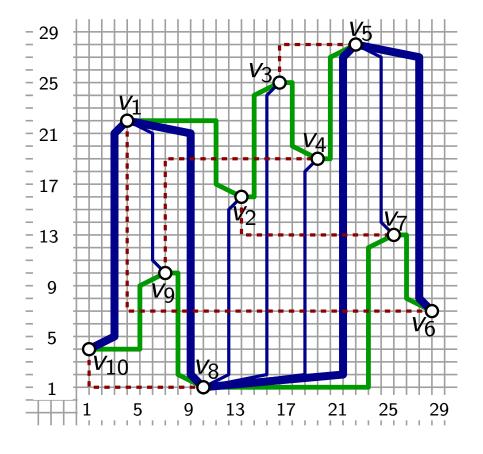


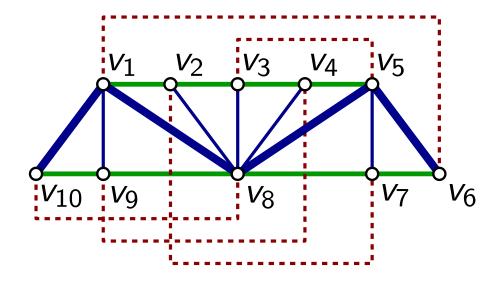




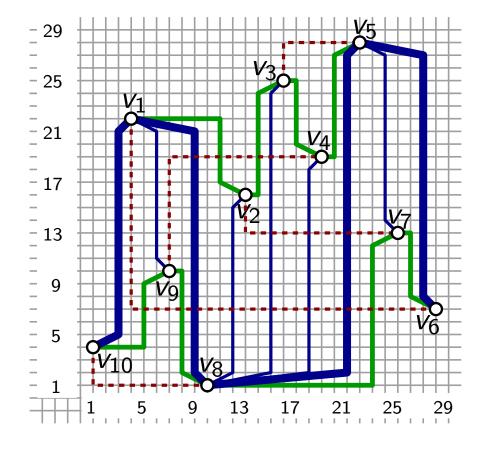








Bends:  $2 \times 1$ Grid size:  $(3n-2)^2$ 

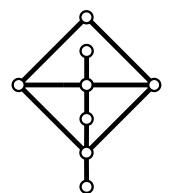


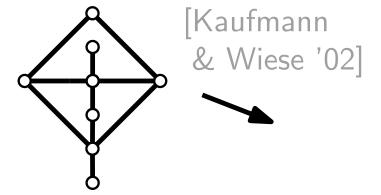
### Overview

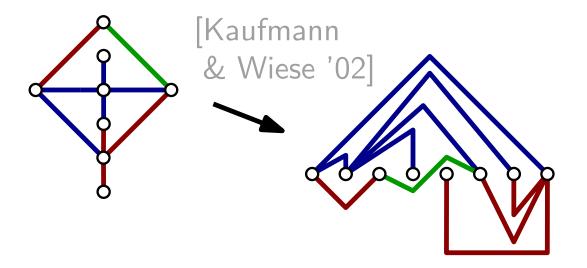
#### Graph classes

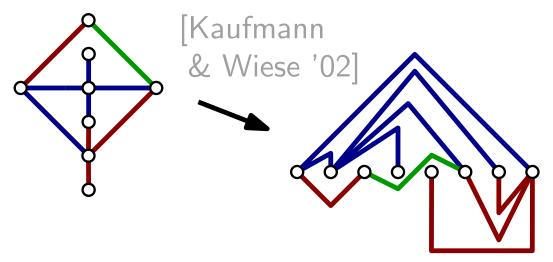
#### Number of bends

Cycle	X	Cycle	$1 \times 1$
Caterpillar	×	Cycle	1  imes 1
Four Matchings			1  imes 1  imes 1  imes 1
Tree	×	Matching	$1 \times 0$
Wheel	×	Matching	$2 \times 0$
Outerpath	×	Matching	$2 \times 1$
Outerplanar	×	Outerplanar	$3 \times 3$
2-page book emb.	×	2-page book emb.	$4 \times 4$
Planar	X	Planar	$6 \times 6$

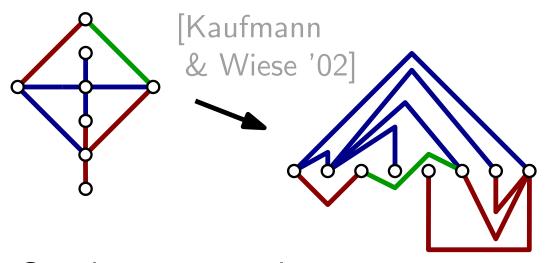




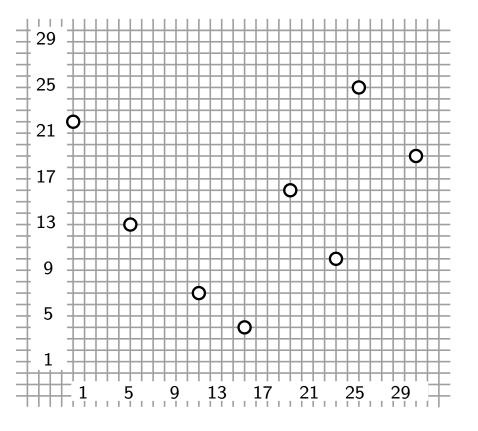


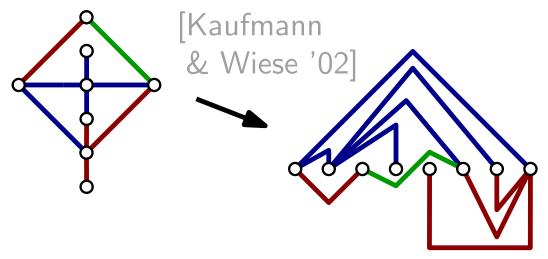


Graph 1: x-coordinates

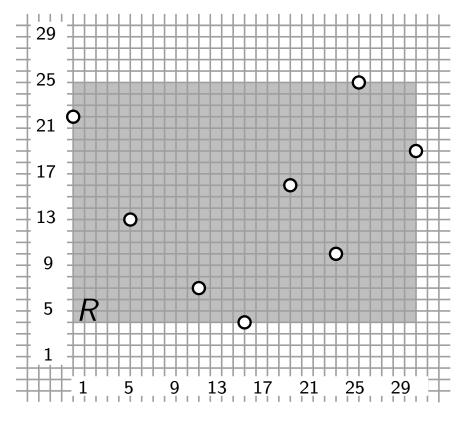


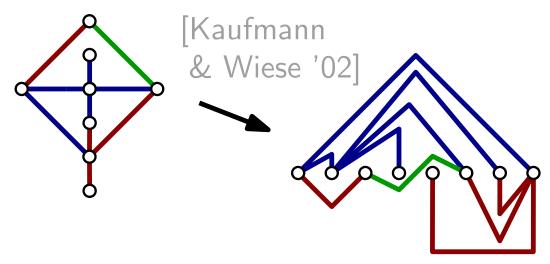
Graph 1: *x*-coordinates





Graph 1: x-coordinates

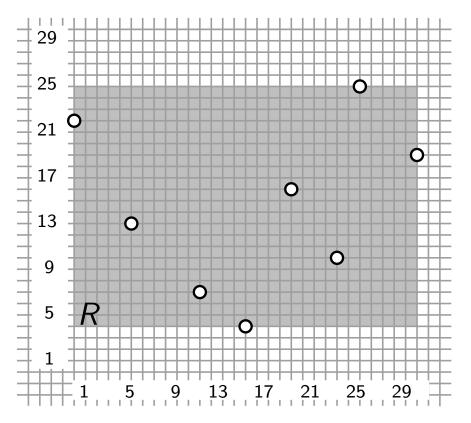


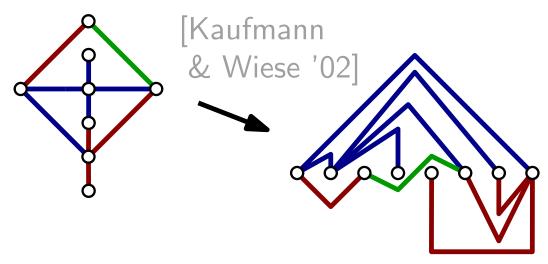


Graph 1: *x*-coordinates

In R: All segments vertical

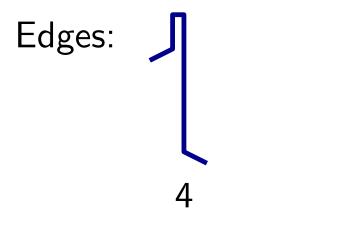
or slanted of *y*-length 1.

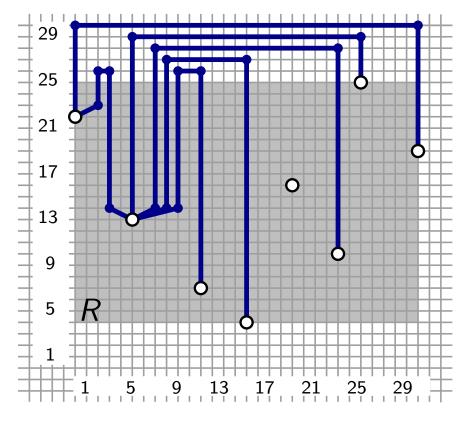


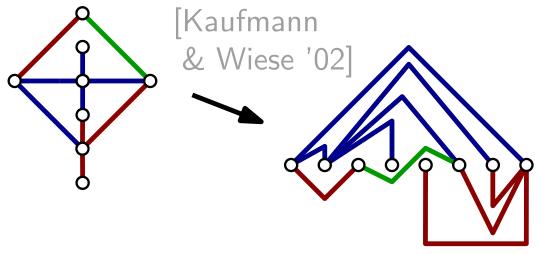


Graph 1: *x*-coordinates

In R: All segments vertical or slanted of y-length 1.







New idea:

Place turns outside of R!

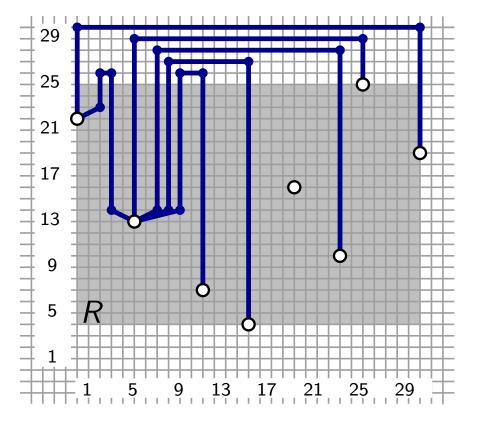
Graph 1: *x*-coordinates

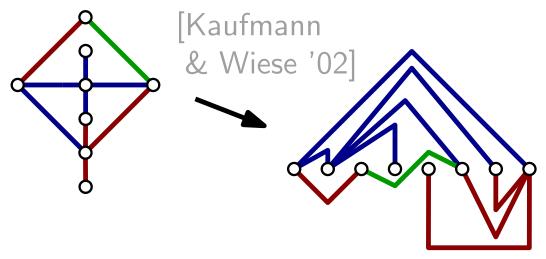
In R: All segments vertical or slanted of y-length 1.

Graph 2: *y*-coordinates

Edges:

4





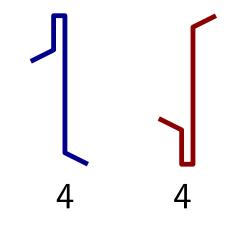
New idea:

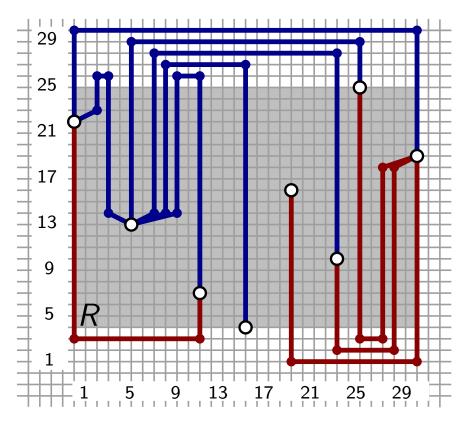
Place turns outside of R!

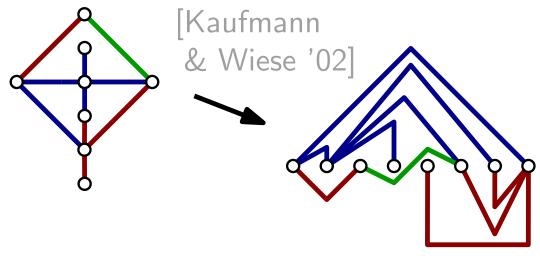
Graph 1: *x*-coordinates

In R: All segments vertical or slanted of y-length 1.

Graph 2: *y*-coordinates







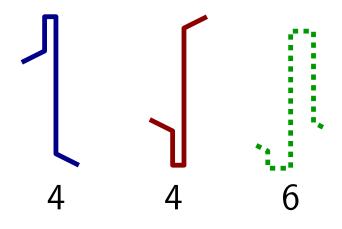
New idea:

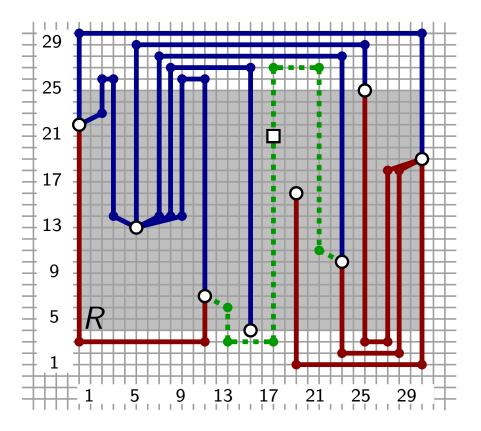
Place turns outside of R!

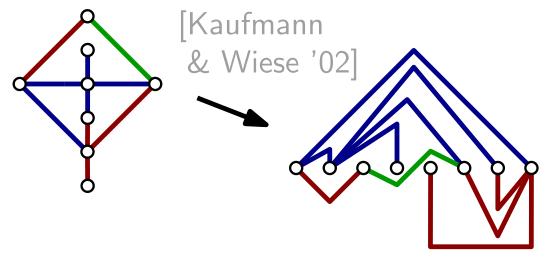
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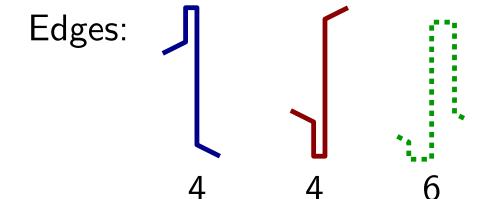


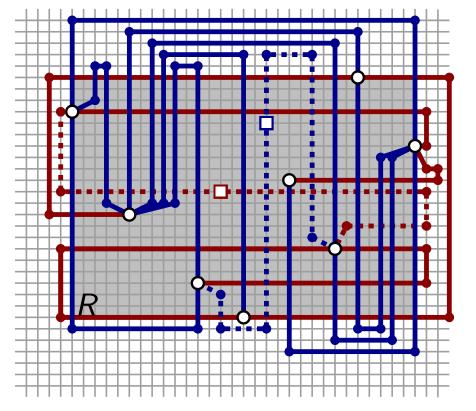
New idea:

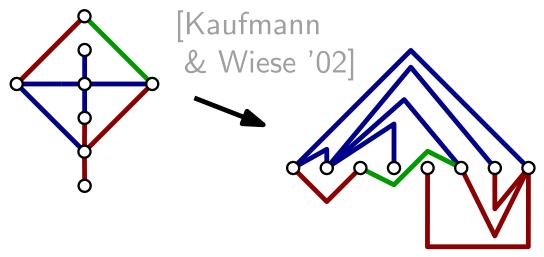
Place turns outside of R!

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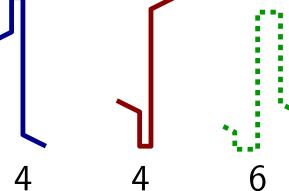


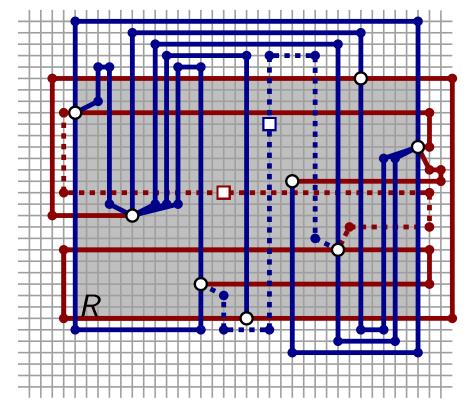
Bends:  $6 \times 6$ Grid size:  $(14n - 26)^2$ 

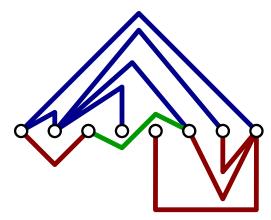
Graph 1: *x*-coordinates

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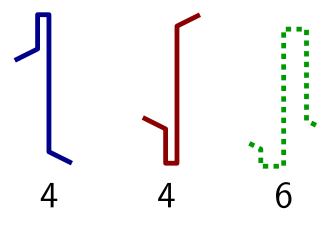


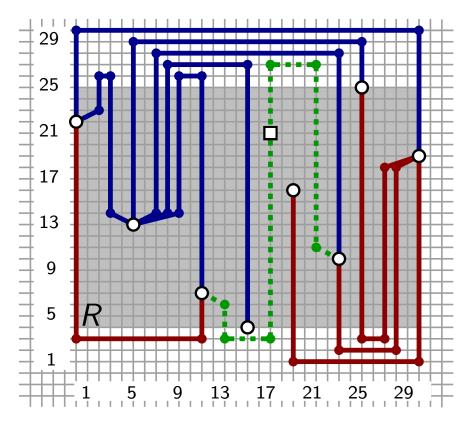


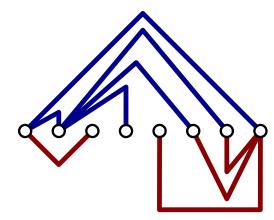
#### Graph 1: *x*-coordinates

In R: All segments vertical or slanted of y-length 1.

Graph 2: y-coordinates



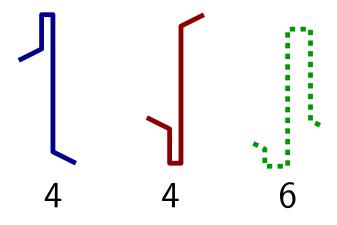


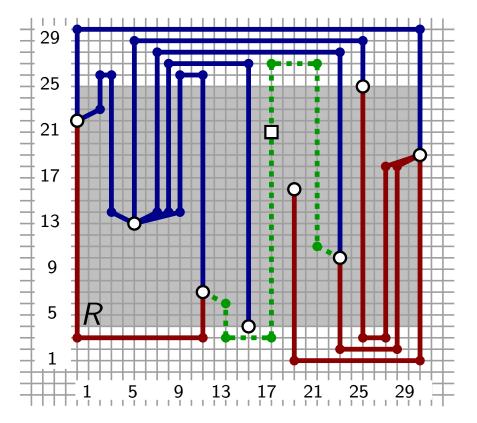


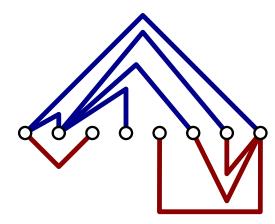
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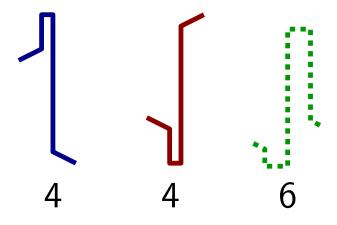


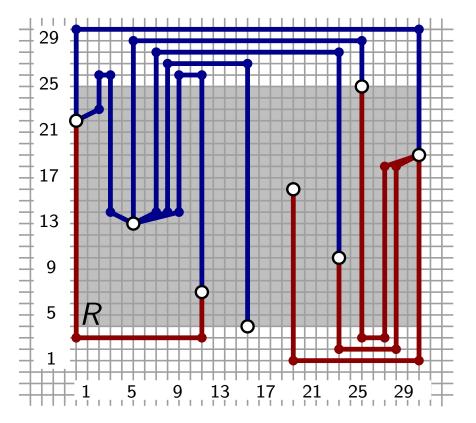


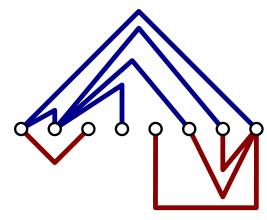
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In R: All segments vertical or slanted of y-length 1.

Graph 2: y-coordinates



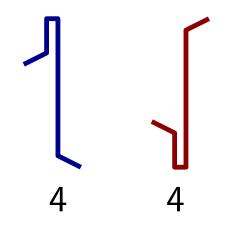


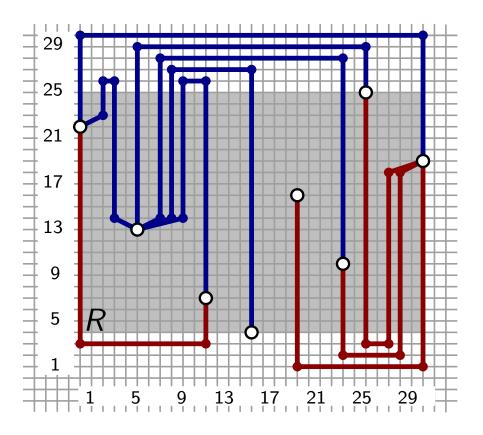


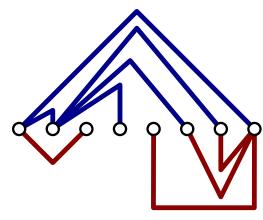
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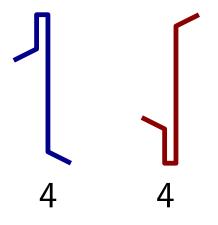


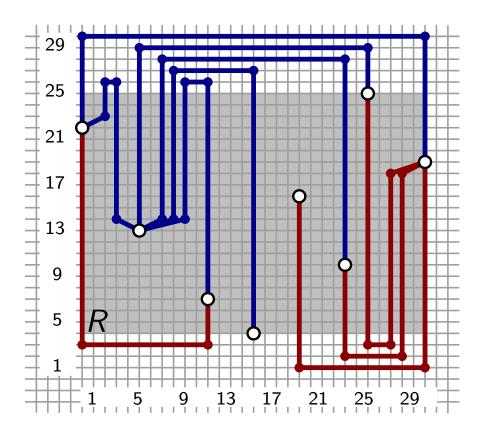
Bends:  $4 \times 4$ Grid size:  $(11n - 32)^2$ 

Graph 1: *x*-coordinates

In R: All segments vertical or slanted of y-length 1.

Graph 2: y-coordinates



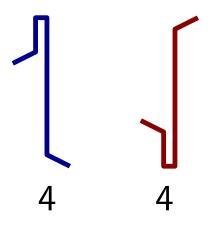


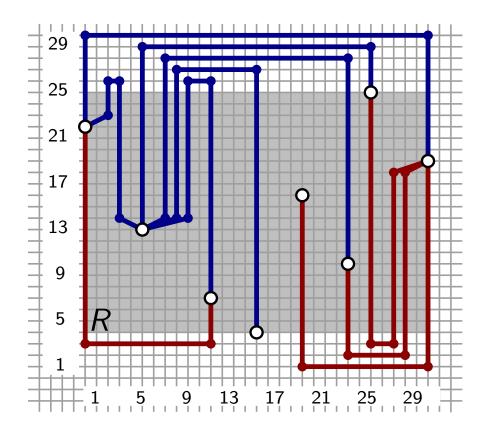
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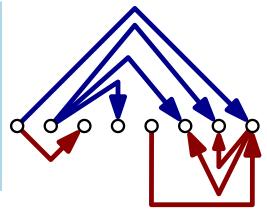


#### Outerplanar × Outerplanar

Decompose into two forests...

[Nash-Williams '64] do o

and direct them!

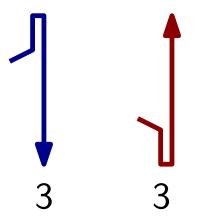


Graph 1: x-coordinates

In R: All segments vertical or slanted of y-length 1.

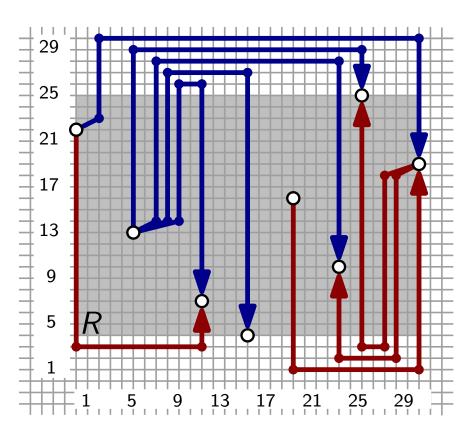
Graph 2: y-coordinates

Edges:



Bends:  $3 \times 3$ 

Grid size:  $(7n - 10)^2$ 



#### Outerplanar × Outerplanar

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[Nash-Williams '64]

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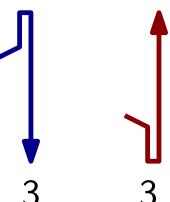
Grid size:  $(7n - 10)^2$ 

Graph 1: x-coordinates

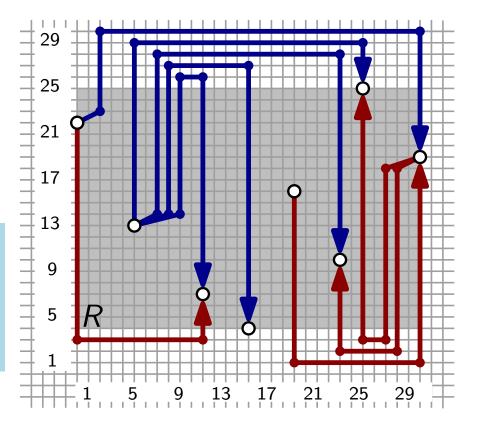
In R: All segments vertical or slanted of y-length 1.

Graph 2: y-coordinates

Edges:



Every vertex has  $\leq 1$  incoming edge from above and  $\leq 1$  from below.



#### Conclusions

#### Graph classes

#### Number of bends

Cycle	×	Cycle	$1 \times 1$	
Caterpillar	×	Cycle	1  imes 1	
Four Matchings			1  imes 1  imes 1  imes 1	
Tree	×	Matching	$1 \times 0$	
Wheel	×	Matching	$2 \times 0$	
Outerpath	×	Matching	$2 \times 1$	
Outerplanar	×	Outerplanar	$3 \times 3$	$\sqrt{}$
2-page book emb.	×	2-page book emb.	$4 \times 4$	
Planar	X	Planar	$6 \times 6$	$\sqrt{}$
				-



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2-page book emb.	×	2-page book emb.	$4 \times 4$
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- All graphs are drawn on the  $O(n) \times O(n)$ -grid.
- $\circ$  All algorithms run in O(n) time.

Reduce bend numbers!

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- Relax constraints on crossing resolution
  - $\Rightarrow$  LacSim

(Simultaneous Embedding with Large-Angle Crossings)

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 Draw each graph RAC (but ignore crossing angles).

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## Our Results (Journal Version)

Graph classes				Number of bends	RAC- SEFE
	Cycle	×	Cycle	1  imes 1	
	Caterpillar	×	Cycle	1  imes 1	
	Four Matchings			1  imes 1  imes 1  imes 1	·
	Tree	×	Matching	$1 \times 0$	
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	2-page book emb.	×	2-page book emb.	$4 \times 4$	
	Planar	×	Planar	$6 \times 6$	

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