To show a statement is truth-functionally true, start proof with no premises and end with the statement:

 $(F \lor G) \to (\neg F \to G)$

(1)

We are aiming for a conditional, so we set up a subproof to get it by → Intro.

Now what we are aiming for is a single

letter, so we look to see how we can get it from earlier steps.

- A. Can we get it by an easy rule? (\land Elim, \rightarrow Elim, \leftrightarrow Elim)
- B. Can we get it from any other complex sentence containing 'G', or if there are non with that, with '¬G'?

Here we find a 'G' in line 1, which is a disjunction. The only thing we can do with that is use vElim, so we set that up.

We see that it is easy to get from 'G' to 'G', so the only work left is to get from 'F' to 'G'.

> To get 'G' under 'F', we notice that we already have ' \neg F' and 'F' available in this subproof. So we can use \perp Intro and then \perp Elim 'G'.

STRATEGY to show that $(F \lor G) \rightarrow (\neg F \rightarrow G)$ is a tautology (is truth-functionally true).

3 Same strategy as in step 2:

$$\begin{vmatrix}
1 \\
2 \\
3 \\
 \end{vmatrix} \begin{vmatrix}
(F \lor G) \\
\neg F \\
 \vdots \\
 \end{vmatrix}$$

$$\begin{vmatrix}
G \\
\neg F \\
\neg G \\
(F \lor G) \\
 \neg F \\
 \neg G \\
 (F \lor G) \\
 \neg F \\
 \neg G \\
 \neg F \\
 \neg F \\
 \neg G \\
 \neg F \\
 \neg F \\
 \neg G \\
 \neg F \\
 \neg F \\
 \neg G \\
 \neg F \\
 \neg F$$





STRATEGYShow that ($\neg A \lor B$) and $A \rightarrow B$ are tautologically equivalent.1 $ \neg A \lor B$	B Derive biconditional from the null (empty) se of premises. As usual, plan to build biconditional by ↔Intro.	et $ \underline{A} \rightarrow \underline{B} $
$ \begin{array}{c} 2\\3\\\\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\$	In the first subproof, build usual way, by ⇔Intro. With this subproof, get 'B' by vH ⇒ 'B' is already in the second subproof. In the first, use the standard strategy we apply when we already have or can easily get a contradiction: use ⊥Intro and ⊥Elim.	nin Elim. $\begin{array}{c c}1\\2\\3\\4\\5\\\end{array} & \begin{vmatrix} \neg A \lor B\\ \hline A\\ \hline A\\3\\4\\5\\\end{vmatrix} & \begin{vmatrix} \neg A\\ \hline A\\ \hline A\\3\\B\\ \hline A\\B\\ \hline A\\1\\B\\ \hline A\\1\\Elim:4\\\end{array}$
	d the disjunction 20	$\begin{vmatrix} \underline{A} \rightarrow \underline{B} \\ \neg A \lor B \\ (\neg A \lor B) \Leftrightarrow (A \rightarrow B) \end{vmatrix}$ $\begin{vmatrix} \underline{A} \\ \neg A \lor B \\ (\neg A \lor B) \Leftrightarrow (A \rightarrow B) \end{vmatrix}$ $\begin{vmatrix} \underline{A} \\ \neg A \lor B \\ \bot \text{ Intro: } 3,4 \\ \underline{B} \\ \bot \text{ Elim: } 5 \\ \underline{B} \\ B \\ \neg A \lor B \\ \neg \text{ Intro: } 3,4 \\ \neg B \\ \neg A \lor B \\ \neg A \lor B \\ \neg A \lor B \\ \neg B \\ \neg A \lor B \\ \neg B \\$

To show that a set is tautologically inconsistent, derive a contradiction from the members of the set.

STRATEGY Show that this set is tautologically inconsistent: $\{A \land (B \leftrightarrow C), A \rightarrow B, (A \land B) \rightarrow D, A \rightarrow (C \rightarrow \neg D)\}$

(I)	
1	$A \land (B \Leftrightarrow C)$	Р
2	A→B	Р
3	$(A \land B) \rightarrow D$	Р
4	$A \land (B \leftrightarrow C) A \rightarrow B (A \land B) \rightarrow D A \rightarrow (C \rightarrow \neg D)$	Р
	· · · · · · · · · · · · · · · · · · ·	
	L	
	1	

 $1 | A \land (B \leftrightarrow C)$ Ρ 2 $A \rightarrow B$ Р 3 $(A \land B) \rightarrow D$ Ρ 4 Р $A \rightarrow (C \rightarrow \neg D)$ 5 \wedge Elim: 1 Α 6 $C \rightarrow \neg D$ \rightarrow Elim: 4,5 ¬D D \bot ⊥ Intro

2 The only '¬' anywhere in this set is in '¬D' inside line 4. So we'll aim for '¬D' and for 'D', then apply ⊥ Intro.

(3) We just saw that '¬D' comes from line 4. To start breaking line 4 apart, we need 'A'. We notice that we can get it from 1 by ∧Elim.

(ৰ)⇒

We also need 'C'. That, too will have to come from breaking apart 1, but not directly. First we pull out 'B \Leftrightarrow C'.

1 2 3 4	$A \land (B \Leftrightarrow C) A \to B (A \land B) \to D A \to (C \to \neg)$	Р Р Р <u>D)</u> Р)
	¬D		
	D ⊥	⊥ Intro)

 $1 \mid A \land (B \leftrightarrow C)$ Р Р 2 $A \rightarrow B$ 3 $(A \land B) \rightarrow D$ Р 4 $A \rightarrow (C \rightarrow \neg D)$ Р 5 ∧ Elim: 1 А 6 $C \rightarrow \neg D$ \rightarrow Elim: 4.5 $7 \mid B \leftrightarrow C$ ∧ Elim: 1 С ¬D D \bot \perp Intro

All that is missing to get 'C' is 'B'. We can get that from 2 and 5 by \rightarrow Elim.

1 2 3 4 5 6 7 8 9 10	$A \land (B \Leftrightarrow A \rightarrow B) (A \land B) \rightarrow A \rightarrow (C \rightarrow A) (C \rightarrow A) (C \rightarrow B) (C \rightarrow C) (C \rightarrow C$	D	1,5 : 1 2,5 7,8
	D ⊥	⊥ Int	ro

(6)=

Now we have all we need to justify ' \neg D'. Next we work on getting 'D'. It must come from line 3. To get it, we must first have 'A \land B', the antecedent of the conditional. We already have both conjuncts separately, so we can build 'A \land B' by \land Intro. That gives us all the steps we need, so all that's left is completing the justifications.

1	A ∧ (B ←	⇒ C) P
2	A→B	P
3	(A ^ B) -	→D P
4	<u>À</u> → (Ć ·	$\rightarrow \neg D)$ P
5	A	∧ Elim: 1
6	$C \rightarrow \neg D$	→Elim: 4,5
7	B ⇔ C	∧ Elim: 1
8	В	→Elim: 2,5
9	С	↔ Elim: 7,8
10	¬D	→Elim: 6,9
11	Α∧Β	∧ Intro: 5,8
12	D	→Elim: 3,11
13	\perp	⊥ Intro: 10,12

1	is tautolog	hat a set of sentences fically inconsistent, ontradiction from the of the set.	{J v K	$J \rightarrow (L \land \neg N)$	how that $K \Leftrightarrow (N \land \neg L)$ ically inconsistent	
3 K ←	K · (L ∧ ¬N) · (N ∧¬L) <u>v N</u>)	P P P	yo easi of t 'L s use	u already have ly get as one n he contradicto Here, if we v N' to contra- entence on lin the common s pplying vElim disjunction	entence 3 K e or can 4 - member 4 - ory pair. 5 e try for - dict the - he 4, we - strategy - n to one -	$ V K \rightarrow (L \land \neg N) \\ (L \lor N) \\ (L \lor N) \\ J \\ L \lor N \\ K \\ L \lor N \\ K \\ V N $
$\begin{array}{c c} 3 & K \\ 4 & \underline{\neg}(L) \\ 5 & J \\ 6 & L \\ & \\ & \\ & \\ & \\ & \\ & \\ & \\ & \\ &$	$(L \land \neg N) \Rightarrow (N \land \neg L) $	P that 'J in line Elim:2,5 disjun the co apply To g		notice cedent irst ' is in we	T NT	
	1 2 3 4 5 6 7 8	L /	P P P Elim:2,5 Elim: 6 Elim: 7	subderiv Apply ↔ separate	ion of the secon ation is similar. Elim to 3, then one of the desir s by AElim.	1
	9 10 11 12 13 14	$\begin{vmatrix} N & & A \\ L & V & V \\ L & V & V \\ \end{bmatrix}$	Elim: 3,9 Elim: 10 Intro: 11 5-8,9-12 tro: 4,13			

To show that a set of (1)Strategy: sentences is tautologically Show that this set is tautologically inconsistent: inconsistent, derive a $\{\neg (P \land Q), R \rightarrow (P \leftrightarrow S), S \land R, S \rightarrow (\neg Q \rightarrow \neg R)\}$ contradiction from the set's members. 1 $\neg (P \land Q)$ 1 $\neg(P \land Q)$ 2) 2 $R \rightarrow (P \leftrightarrow S)$ 2 $R \rightarrow (P \leftrightarrow S)$ Look for a negation 3 3 $S \wedge R$ $S \wedge R$ we already have as one member of the 4 4 $S \rightarrow (\neg Q \rightarrow \neg R)$ $S \rightarrow (\neg Q \rightarrow \neg R)$ contradictory pair. Here, try for 'P \land Q' to contradict the sentence on line 1. $P \land Q$ \bot \bot $\neg (P \land O)$ 1 3 2 3 To get ' $P \land Q$ ', we need $R \rightarrow (P \leftrightarrow S)$ $\neg (P \land Q)$ 1 $S \wedge R$ each conjunct. First get 'P'. 'P' 2 $R \rightarrow (P \leftrightarrow S)$ 4 $S \rightarrow (\neg Q \rightarrow \neg R)$ occurs only in lines 1 and 2. We 3 $S \wedge R$ can't break apart line 1, so we must 4 $\underline{S \rightarrow (\neg Q \rightarrow \neg R)}$ $P \leftrightarrow S$ use line 2. We need to separate out 5 R ∧Elim: (its consequent. $6 | P \leftrightarrow S$ \rightarrow Elim: 2, Р Р 4 Q P ∧ Q To get 'P \Leftrightarrow S' from 2, we must Q \bot P∧Q first get the antecedent, 'R'. Get 'R' from 3 by $\wedge E$. \bot 5 6, 'Q' appears above only 7) ¬Q' leads easily to To get 'P' from in lines 1 & 4. We can't get $^{\prime} \neg R'$. With $^{\prime}R'$ 'P \Leftrightarrow S', we need 'S', 'Q' directly, so plan to get a which we can also get (line 5), we have a contradiction. 'R' and ' \neg R' from 3. look promising. Assume '¬Q' contradiction. This to help us get $\neg R'$ and ' \perp '. completes the proof. \neg (P \land Q) 1 2 $R \rightarrow (P \leftrightarrow S)$ $1 | \neg (P \land Q)$ 1 $\neg (P \land Q)$ 3 $S \wedge R$ 2 $R \rightarrow (P \leftrightarrow S)$ 2 $R \rightarrow (P \leftrightarrow S)$ 3 4 S ¬R) 3 $S \wedge R$ $S \wedge R$ 5 R \wedge Elim: 3 4 4 $S \rightarrow (\neg Q \rightarrow \neg R)$ $\underline{S \rightarrow (\neg Q \rightarrow \neg R)}$ 5 6 $P \leftrightarrow S$ \rightarrow Elim: 2,5 5 6 7 8 R R ∧ Elim: 3 \land Elim: 3 7 S \wedge Elim: 3 6 $P \leftrightarrow S$ \rightarrow Elim: 2,5 $P \leftrightarrow S$ \rightarrow Elim: 2,5 8 P 7 ↔Elim: 6,7 S S ∧ Elim: 3 \wedge Elim: 3 8 Ρ Р \leftrightarrow Elim: 6,7 ↔ Elim: 6,7 $\begin{array}{c} Q \\ P \land Q \end{array}$ 9 9 <u>¬Q</u> 10 ¬R →Elim: 4,7 \bot ¬R ¬R →Elim: 9,10 11 \bot 12 \perp Intro: 5,11 $\neg \neg O$ 13 $\neg \neg O$ ¬ Intro: 9-12 Q 14 Q ¬¬Elim: 13 P`∧Q 15 $P \land Q$ ∧Intro: 8,14 16 ⊥ \perp Intro: 1,15